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Strong converse theorems using Rényi entropies

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Abstract

We use a Rényi entropy method to prove strong converse theorems for certain information-theoretic tasks which involve local operations and quantum or classical communication between two parties. These include state redistribution, coherent state merging, quantum state splitting, measurement compression with quantum side information, randomness extraction against quantum side information, and data compression with quantum side information. The method we employ in proving these results extends ideas developed by Sharma [41], which he used to give a new proof of the strong converse theorem for state merging. For state redistribution, we prove the strong converse property for the boundary of the entire achievable rate region in the (e, q) -plane, where e and q denote the entanglement cost and quantum communication cost, respectively. In the case of measurement compression with quantum side information, we prove a strong converse theorem for the classical communication cost, which is a new result extending the previously known weak converse. For the remaining tasks, we provide new proofs for strong converse theorems previously established using smooth entropies. For each task, we obtain the strong converse theorem from explicit bounds on the figure of merit of the task in terms of a Rényi generalization of the optimal rate. Hence, we identify candidates for the strong converse exponents for each task discussed in this paper. To prove our results, we establish various new entropic inequalities, which might be of independent interest. These involve conditional entropies and mutual information derived from the sandwiched Rényi divergence. In particular, we obtain novel bounds relating these quantities, as well as the Rényi conditional mutual information, to the fidelity of two quantum states.

1 Introduction

1.1 Strong converse theorems and the Rényi entropy method

One of the primary goals of quantum information theory is to find optimal rates of information-theoretic tasks, such as data compression, information transmission through a noisy quantum channel, and entanglement manipulation. Depending on the specific task in question, the optimal rate is either an optimal *cost*, quantifying the minimum rate at which an available resource is consumed in accomplishing the task, or an optimal *gain*, quantifying the maximum rate at which a desired target resource is produced in the process. For any rate above (below) the optimal cost (gain) there is a corresponding protocol for accomplishing the task successfully. That is, the error, ε_n , incurred in the protocol for n uses of the underlying resource vanishes in

the asymptotic limit ($n \rightarrow \infty$).¹ Such rates are said to be *achievable*. In contrast, for protocols with non-achievable rates, the error does not vanish asymptotically. The optimal rate of an information-theoretic task is said to satisfy the *strong converse property* if for any sequence of protocols with a non-achievable rate the error ε_n is not only bounded away from zero but necessarily converges to one in the asymptotic limit. Moreover, this convergence is exponential² in n , that is,

$$\varepsilon_n \geq 1 - \exp(-Kn), \quad (1.1)$$

for some positive constant K . A *strong converse theorem* establishes the strong converse property of the optimal rate of an information-theoretic task, and hence serves to identify the latter as a sharp rate threshold for the task.

For information transmission through classical noisy channels, the strong converse theorem was first proved by Wolfowitz [53]. An alternate proof of this theorem was later given by Arimoto [2] by employing the properties of a quantity which is sometimes referred to as the Gallager function [35]. Ogawa and Nagaoka [32] extended this method to the quantum setting to prove the strong converse property of the capacity of a classical-quantum channel, which was also proved concurrently by Winter [50] using the method of types. Nagaoka [31] further developed Arimoto's idea to give a new proof of this result. To this end, he employed a Rényi divergence and its monotonicity under completely positive, trace-preserving maps (also called the *data processing inequality*), establishing what we refer to as the 'Rényi entropy method' in this paper. Later, Polyanskiy and Verdú [35] realized that it is possible to establish converse bounds by employing any divergence satisfying the data processing inequality. In [18] the Rényi entropy method was used to obtain strong converse theorems in network information theory.

Different quantum generalizations of the α -Rényi divergence have been introduced [34, 30, 48] and their monotonicity under quantum operations for certain ranges of the Rényi parameter α has been established [34, 19, 3]. The Rényi entropy method has since been successfully employed to prove strong converse theorems for classical channel coding with entangled inputs for a large class of quantum channels with additive Holevo capacity [26]. More recently, strong converse theorems were proved for classical information transmission through entanglement-breaking and Hadamard channels [48] and quantum information transmission through generalized dephasing channels [46]. For the task of quantum hypothesis testing, Mosonyi and Ogawa [29] established the 'sandwiched Rényi divergence' of order α [30, 48] as a strong converse exponent. This was generalized by Cooney *et al.* [11], who established the sandwiched Rényi divergence as a strong converse exponent in adaptive channel discrimination for certain channels. Hayashi and Tomamichel [24] showed that, in binary quantum hypothesis testing, for a composite alternative hypothesis the strong converse exponents are given by a Rényi mutual information and Rényi conditional entropy defined in terms of the sandwiched Rényi divergence.

Application of the Rényi entropy method to prove the strong converse property for an information-theoretic task involving local operations and classical communication (LOCC) between two parties (say, Alice and Bob) was considered by Hayashi *et al.* [23] in the context of entanglement concentration (see also [22]). More recently, Sharma [41] used the Rényi entropy method to establish the strong converse theorem for the task of *state merging*: Alice and Bob initially share a bipartite state and the aim is for Alice to transfer her part of the state to Bob by sending information to him through a noiseless classical channel. Both Alice and Bob are also

¹For example, in the case of information transmission through a memoryless noisy channel, n denotes the number of independent uses of the channel.

²The requirement of exponential convergence is sometimes relaxed when defining the strong converse property.

allowed to make use of prior shared entanglement between them, to assist them in achieving this task. In this case monotonicity of a relevant Rényi divergence under LOCC plays a pivotal role in establishing the strong converse for the optimal entanglement cost.³

In this paper, we further extend the Rényi entropy method to establish strong converse theorems for a range of information-theoretic tasks which involve local operations and quantum or classical communication (LOQC or LOCC) between two parties. These tasks (which are elucidated at the beginning of each section respectively) include state redistribution (with and without feedback), coherent state merging, quantum state splitting, measurement compression with quantum side information, randomness extraction against quantum side information, and data compression with quantum side information. Strong converse theorems for some of these tasks were previously obtained in the so-called smooth entropy framework introduced by Renner [38]. This was done for the quantum communication cost of state redistribution by Berta *et al.* [8], and for randomness extraction against quantum side information and data compression with quantum side information by Tomamichel [42].

Two inequalities which we use in proving the strong converse theorems are due to van Dam and Hayden [55]: the first bounds the fidelity between two states in terms of the difference of their Rényi entropies, and the second is a certain subadditivity property of the α -Rényi entropy of a bipartite state. In addition, we establish various new inequalities involving conditional entropies and mutual information derived from the sandwiched Rényi divergence of order α [30, 48]. These inequalities play a crucial role in the proofs of our strong converse theorems and might also be of independent interest.

Let us use the example of quantum data compression to explain the key step of the Rényi entropy method which establishes the strong converse property (cf. [21]). Let ρ denote the source state of a discrete, memoryless source. In this case, the optimal rate r^* is the minimum rate of data compression and is given [40] by the von Neumann entropy of the source: $r^* = S(\rho) := -\text{Tr}(\rho \log \rho)$.⁴ Consider a data compression protocol of rate r . The key step of the Rényi entropy method applied to quantum data compression is to prove that for values of a real parameter $\alpha > 1$, there exists a positive constant $\kappa(\alpha)$ such that the error, ε_n , incurred for n independent uses of the source, satisfies the following bound [41]:

$$\varepsilon_n \geq 1 - \exp \{-n\kappa(\alpha) [S_\alpha(\rho) - r]\}. \quad (1.2)$$

Here, $\kappa(\alpha) = (\alpha - 1)/(2\alpha)$, and $S_\alpha(\rho)$ is the Rényi entropy of the source state ρ :

$$S_\alpha(\rho) := \frac{1}{1 - \alpha} \log \text{Tr} \rho^\alpha,$$

satisfying (cf. Proposition 2.5)

$$S_\alpha(\rho) \leq S_{\alpha'}(\rho) \leq S(\rho) \quad \text{for } \alpha \geq \alpha' > 1 \quad \text{and} \quad \lim_{\alpha \rightarrow 1} S_\alpha(\rho) = S(\rho) = r^*. \quad (1.3)$$

The inequality (1.2), along with the statements in (1.3), readily imply the statement (1.1) of the strong converse theorem. To see this, suppose $r < r^* = S(\rho)$. Then there is a $\delta > 0$ such that $r + \delta < r^*$. Moreover, (1.3) implies that for every $\delta > 0$ there is an $\alpha_0 > 1$ such that $S_{\alpha_0}(\rho) > r^* - \delta$. Combining the two inequalities, we get $r < r^* - \delta < S_{\alpha_0}(\rho)$, and inserting this in (1.2) yields the strong converse condition as stated in (1.1) with the choice $K := \kappa(\alpha_0)[S_{\alpha_0}(\rho) - r] > 0$.

³This result also follows from prior work by various authors [52, 6, 42].

⁴In this paper all logarithms and exponentials are taken to base 2.

1.2 Main results and structure of the paper

In the present paper, we prove strong converse theorems for the following tasks using the Rényi entropy method:

- (a) state redistribution (with and without quantum feedback), coherent state merging and quantum state splitting;
- (b) measurement compression with quantum side information (QSI);
- (c) randomness extraction against QSI;
- (d) data compression with QSI.

Previously, Tomamichel [42] proved strong converse theorems for randomness extraction against QSI and data compression with QSI using the *smooth entropy framework*. Recently, Berta *et al.* [8] proved a strong converse theorem for the quantum communication cost in state redistribution (which holds even in the presence of quantum feedback). However, their strong converse theorem did not establish the strong converse property for the boundary of the entire achievable rate region in the (e, q) -plane, where e and q denote the entanglement cost and quantum communication cost, respectively (see Figure 1). In this paper, we fill this gap with Theorem 3.2 (for the case of no feedback) and Theorem 3.4 (for the case of feedback). The study of the strong converse for state redistribution with feedback was inspired by [8], where this protocol was first defined. Note that, following discussions with the authors of the present paper, Berta *et al.* have now also obtained a strong converse theorem for the entire achievable region [9] for state redistribution with and without feedback. In the case of measurement compression with quantum side information, our strong converse theorem for the classical communication cost is a new result, which strengthens the previously established weak converse of [47]. We also provide alternative proofs for the strong converse theorems of the protocols in items **c** and **d** in the above list using the Rényi entropy approach.

Our strong converse theorems are given in terms of various Rényi generalizations (see Section 2.3) of the optimal rates of the protocols, whose properties we derive in Section 2.4. In particular, we establish various new inequalities involving conditional entropies and mutual information derived from the sandwiched Rényi divergence. These include novel bounds relating these entropic quantities, as well as the Rényi conditional mutual information (defined in 2.14), to the fidelity of two quantum states (see Proposition 2.8). These bounds play a key role in our proofs.

This paper is structured as follows. In Section 2 we set the notation, introduce definitions of the required Rényi entropic quantities, and discuss their properties. In the following sections we subsequently prove strong converse theorems for the protocols of state redistribution (Section 3), measurement compression with quantum side information (Section 4), randomness extraction against quantum side information (Section 5), and data compression with quantum side information (Section 6). In the case of state redistribution, we elucidate the fact that it serves as a generalization of coherent state merging (Section 3.5) and quantum state splitting (Section 3.6), proving strong converse theorems for the latter protocols as well. We summarize our results and discuss open questions in Section 7. Appendices A and B contain some of the proofs of Sections 2.4 and 3.4.

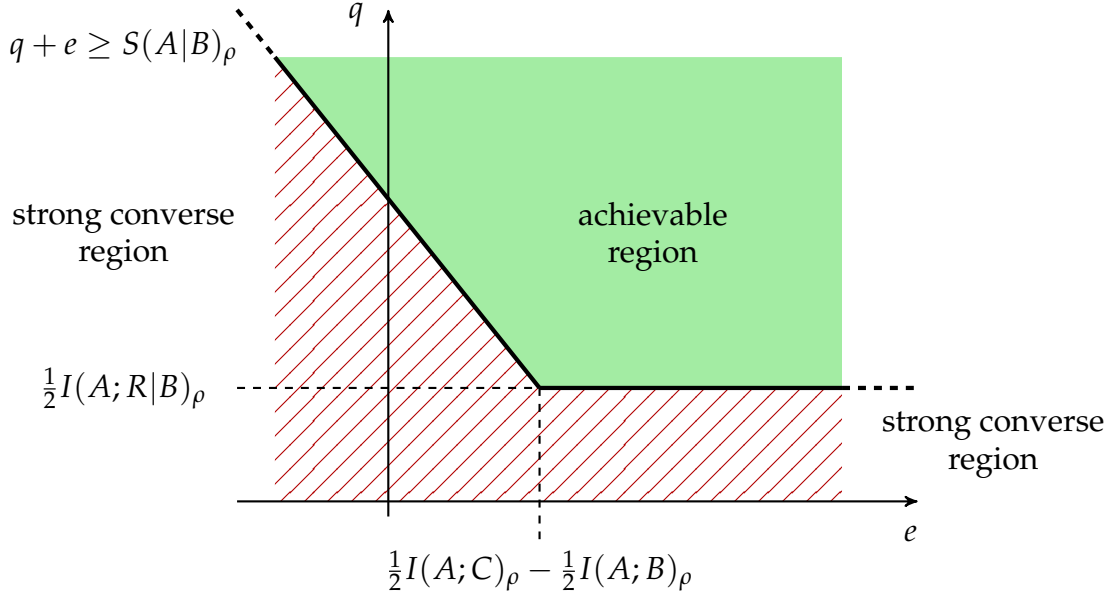


Figure 1: Plot of the plane of rate pairs (e, q) for state redistribution, where e is the entanglement cost (3.5) and q is the quantum communication cost (3.4). The shaded area is the region of achievable rate pairs defined by $\{(e, q) : q + e \geq S(A|B)_\rho \text{ and } q \geq \frac{1}{2}I(A; R|B)_\rho\}$. The hatched area is the strong converse region, as proved in the present paper (Theorem 3.2) and by Berta *et al.* [9]. For a definition of the state redistribution protocol, see Section 3.1.

2 Preliminaries

2.1 Notation & definitions

For a finite-dimensional Hilbert space \mathcal{H} we denote the set of linear operators acting on \mathcal{H} by $\mathcal{B}(\mathcal{H})$. We define the set of positive semi-definite operators $\mathcal{P}(\mathcal{H}) := \{X \in \mathcal{B}(\mathcal{H}) : X \geq 0\}$ and refer to $P \in \mathcal{P}(\mathcal{H})$ simply as a positive operator. The set $\mathcal{D}(\mathcal{H})$ of density operators (or quantum states) on \mathcal{H} is the set of positive operators with unit trace: $\mathcal{D}(\mathcal{H}) := \{\rho \in \mathcal{P}(\mathcal{H}) : \text{Tr} \rho = 1\}$. The support $\text{supp } Q$ of an operator Q is defined as the orthogonal complement of its kernel. We write $A \not\perp B$ if $\text{supp } A \cap \text{supp } B$ contains at least one non-zero vector. For a quantum system A with associated Hilbert space \mathcal{H}_A we write $|A| := \dim \mathcal{H}_A$. If B is another quantum system with associated Hilbert space \mathcal{H}_B , then we set $\mathcal{H}_{AB} := \mathcal{H}_A \otimes \mathcal{H}_B$. We write $A \cong B$ for quantum systems A and B whose underlying Hilbert spaces are isomorphic.

For a pure state $|\psi_A\rangle \in \mathcal{H}_A$ we make the abbreviation $\psi_A \equiv |\psi\rangle\langle\psi|_A \in \mathcal{D}(\mathcal{H}_A)$. We denote by I_A the identity operator acting on \mathcal{H}_A , and by id_A the identity superoperator acting on $\mathcal{D}(\mathcal{H}_A)$. The completely mixed state on \mathcal{H}_A is denoted by $\pi_A := |A|^{-1}I_A$. Let \mathcal{X} be some finite set and \mathcal{H}_X the associated Hilbert space with orthonormal basis $\{|x\rangle\}_{x \in \mathcal{X}}$. Then the quantum embedding of a classical register X corresponding to \mathcal{X} is defined as the subspace of diagonal matrices in $\mathcal{D}(\mathcal{H}_X)$. A classical-quantum (c-q) state $\rho_{XB} \in \mathcal{D}(\mathcal{H}_{XB})$ is defined as $\rho_{XB} := \sum_{x \in \mathcal{X}} p_x |x\rangle\langle x|_X \otimes \rho_B^x$ where $\rho_B^x \in \mathcal{D}(\mathcal{H}_B)$ for all $x \in \mathcal{X}$. For quantum systems $A \cong B$, a maximally entangled state (MES) $|\Phi_{AB}\rangle$ is defined as

$$|\Phi_{AB}\rangle := \frac{1}{\sqrt{|A|}} \sum_{i=1}^{|A|} |i_A\rangle \otimes |i_B\rangle,$$

where $\{|i_A\rangle\}_{i=1}^{|A|}$ and $\{|i_B\rangle\}_{i=1}^{|B|}$ are orthonormal bases for \mathcal{H}_A and \mathcal{H}_B , respectively. We also use the notation Φ_{AB}^k to explicitly indicate the Schmidt rank $k = |A| = |B|$ of the MES.

A quantum operation Λ is a linear, completely positive, trace-preserving (CPTP) map. For a quantum operation $\Lambda: \mathcal{D}(\mathcal{H}_A) \rightarrow \mathcal{D}(\mathcal{H}_B)$ we write $\Lambda: A \rightarrow B$. For quantum states $\rho, \sigma \in \mathcal{D}(\mathcal{H})$, we define the *fidelity* $F(\rho, \sigma)$ between ρ and σ as

$$F(\rho, \sigma) := \|\sqrt{\rho}\sqrt{\sigma}\|_1.$$

The *von Neumann entropy* of a state $\rho_A \in \mathcal{D}(\mathcal{H}_A)$ is given by $S(\rho_A) := -\text{Tr}(\rho_A \log \rho_A)$, and we use the notation $S(A)_\rho \equiv S(\rho_A)$. The *quantum relative entropy* of a state $\rho \in \mathcal{D}(\mathcal{H})$ and a positive operator $\sigma \in \mathcal{P}(\mathcal{H})$ is defined by

$$D(\rho\|\sigma) := \text{Tr}[\rho(\log \rho - \log \sigma)] \quad (2.1)$$

if $\text{supp } \rho \subseteq \text{supp } \sigma$, and set to be equal to $+\infty$ otherwise. Furthermore, we consider the following quantities for a tripartite state $\rho_{ABC} \in \mathcal{D}(\mathcal{H}_{ABC})$ and its marginals:

- the *quantum conditional entropy* $S(A|B)_\rho := S(AB)_\rho - S(B)_\rho$
- the *quantum mutual information* $I(A; B)_\rho := S(A)_\rho + S(B)_\rho - S(AB)_\rho$
- the *conditional quantum mutual information* $I(A; B|C)_\rho := S(A|C)_\rho + S(B|C)_\rho - S(AB|C)_\rho$

2.2 Schatten norms

Definition 2.1 (Schatten p -(quasi)norm). Let $M \in \mathcal{B}(\mathcal{H})$ and $p > 0$. Then we define

$$\|M\|_p := (\text{Tr} |M|^p)^{1/p},$$

where $|M| := \sqrt{M^\dagger M}$. For $1 \leq p \leq \infty$, the functional $\|\cdot\|_p$ defines a norm, the *Schatten p -norm*.

Theorem 2.2. Let $M, N \in \mathcal{B}(\mathcal{H})$.

- (i) *Hölder's inequality*: Let $1 \leq p \leq \infty$ and q be the Hölder conjugate of p defined by $\frac{1}{p} + \frac{1}{q} = 1$. Then

$$\|MN\|_1 \leq \|M\|_p \|N\|_q. \quad (2.2)$$

- (ii) *McCarthy's inequalities* [28]: For $p \in (0, 1)$ we have

$$\|M + N\|_p^p \leq \|M\|_p^p + \|N\|_p^p, \quad (2.3)$$

whereas for $1 \leq p \leq \infty$ we have

$$\|M\|_p^p + \|N\|_p^p \leq \|M + N\|_p^p. \quad (2.4)$$

Remark 2.3. For $1 \leq p \leq \infty$, the functional $\|\cdot\|_p$ satisfies the triangle inequality, which for the Schatten p -norms is also known as the *Minkowski inequality*:

$$\|M + N\|_p \leq \|M\|_p + \|N\|_p. \quad (2.5)$$

Hence, $\|\cdot\|_p$ defines a norm for this range of p . However, for $p \in (0, 1)$ the Minkowski inequality (2.5) fails to hold, and we have the weaker inequality (2.3) instead. Therefore, $\|\cdot\|_p$ only defines a quasinorm for the range $p \in (0, 1)$.

2.3 Rényi entropies

Definition 2.4 ([30, 48]).

- (i) Let $\alpha \in (0, \infty) \setminus \{1\}$, and $\rho, \sigma \in \mathcal{P}(\mathcal{H})$. If $\text{supp } \rho \subseteq \text{supp } \sigma$ for $\alpha > 1$ or $\rho \not\subseteq \sigma$ for $\alpha \in (0, 1)$, the *sandwiched Rényi divergence of order α* is defined as

$$\begin{aligned}\tilde{D}_\alpha(\rho\|\sigma) &:= \frac{1}{\alpha-1} \log \left[(\text{Tr } \rho)^{-1} \text{Tr} \left\{ \left(\sigma^{(1-\alpha)/2\alpha} \rho \sigma^{(1-\alpha)/2\alpha} \right)^\alpha \right\} \right] \\ &= \frac{2\alpha}{\alpha-1} \log \left\| \rho^{1/2} \sigma^{(1-\alpha)/2\alpha} \right\|_{2\alpha} - \frac{1}{\alpha-1} \log \text{Tr } \rho.\end{aligned}$$

Otherwise, we set $\tilde{D}_\alpha(\rho\|\sigma) = \infty$. Note that for $[\rho, \sigma] = 0$ the sandwiched Rényi divergence reduces to the usual α -relative Rényi entropy $D_\alpha(\rho\|\sigma) := \frac{1}{\alpha-1} \log \left[(\text{Tr } \rho)^{-1} \text{Tr} (\rho^\alpha \sigma^{1-\alpha}) \right]$ (see e.g. [34]).

- (ii) For $\rho \in \mathcal{D}(\mathcal{H})$ and $\alpha \in (0, \infty) \setminus \{1\}$, the *Rényi entropy of order α* is defined as

$$S_\alpha(\rho) := -D_\alpha(\rho\|I) = -\tilde{D}_\alpha(\rho\|I).$$

Note that $S_0(\rho) = \lim_{\alpha \rightarrow 0} S_\alpha(\rho) = \log \text{rk } \rho$, where $\text{rk } \rho$ denotes the rank of ρ . We use the notation $S_\alpha(A)_\rho \equiv S_\alpha(\rho_A)$ for $\rho_A \in \mathcal{D}(\mathcal{H}_A)$.

- (iii) For $\rho_{AB} \in \mathcal{D}(\mathcal{H}_{AB})$ and $\alpha \in (0, \infty) \setminus \{1\}$, the *Rényi conditional entropy of order α* is defined as

$$\tilde{S}_\alpha(A|B)_\rho := -\min_{\sigma_B} \tilde{D}_\alpha(\rho_{AB}\|I_A \otimes \sigma_B),$$

and the *Rényi mutual information of order α* is defined as [20]

$$\tilde{I}_\alpha(A; B)_\rho := \min_{\sigma_B} \tilde{D}_\alpha(\rho_{AB}\|\rho_A \otimes \sigma_B).$$

We have [30, 48]

$$\lim_{\alpha \rightarrow 1} \tilde{D}_\alpha(\rho\|\sigma) = D(\rho\|\sigma). \quad (2.6)$$

Therefore, in the subsequent discussion of the sandwiched Rényi divergence $\tilde{D}_\alpha(\rho\|\sigma)$ and its derived quantities, we will use the full range $\alpha \geq 0$, setting $\tilde{D}_1(\rho\|\sigma) = D(\rho\|\sigma)$ and $\tilde{D}_0(\rho\|\sigma) = \lim_{\alpha \rightarrow 0} \tilde{D}_\alpha(\rho\|\sigma)$ [14]. In the next proposition we collect a few properties of the Rényi quantities defined above:

Proposition 2.5 ([30, 48]). *Let $\rho \in \mathcal{D}(\mathcal{H})$ and $\sigma \in \mathcal{P}(\mathcal{H})$. The sandwiched Rényi divergence and entropies derived from it satisfy the following properties:*

- (i) *Monotonicity: If $0 < \alpha \leq \beta$, then $\tilde{D}_\alpha(\rho\|\sigma) \leq \tilde{D}_\beta(\rho\|\sigma)$.*
- (ii) *Positivity and dimension bound: Let $d = \dim \mathcal{H}$, then we have $0 \leq S_\alpha(\rho) \leq \log d$ for all $\alpha \geq 0$. The extremal values are achieved for pure states and completely mixed states, respectively.*
- (iii) *Additivity: For all $\alpha > 0$ and $\rho_i \in \mathcal{D}(\mathcal{H}_i)$, $\sigma_i \in \mathcal{P}(\mathcal{H}_i)$ for $i = 1, 2$ such that the conditions on their supports given by Definition 2.4 hold, we have*

$$\tilde{D}_\alpha(\rho_1 \otimes \rho_2\|\sigma_1 \otimes \sigma_2) = \tilde{D}_\alpha(\rho_1\|\sigma_1) + \tilde{D}_\alpha(\rho_2\|\sigma_2),$$

$$S_\alpha(\rho_1 \otimes \rho_2) = S_\alpha(\rho_1) + S_\alpha(\rho_2).$$

Furthermore, additivity also holds for the Rényi conditional entropy and mutual information [3, 24]: If ρ_{AB} and $\sigma_{A'B'}$ are bipartite states and $\alpha \geq 1/2$, then

$$\begin{aligned}\tilde{S}_\alpha(AA'|BB')_{\rho \otimes \sigma} &= \tilde{S}_\alpha(A|B)_\rho + \tilde{S}_\alpha(A'|B')_\sigma, \\ \tilde{I}_\alpha(AA'; BB')_{\rho \otimes \sigma} &= \tilde{I}_\alpha(A; B)_\rho + \tilde{I}_\alpha(A'; B')_\sigma.\end{aligned}$$

(iv) *Invariance under isometries:* Let $U: \mathcal{H} \rightarrow \mathcal{H}'$ be an isometry. Then $\tilde{D}_\alpha(U\rho U^\dagger \| U\sigma U^\dagger) = \tilde{D}_\alpha(\rho \| \sigma)$ and $S_\alpha(U\rho U^\dagger) = S_\alpha(\rho)$ for all $\alpha \geq 0$.

(v) *Duality for the Rényi entropy:* Let $|\psi_{AB}\rangle$ be a pure state and set $\rho_A = \text{Tr}_B \psi_{AB}$ and $\rho_B = \text{Tr}_A \psi_{AB}$. Then $S_\alpha(A)_\rho = S_\alpha(B)_\rho$ for all $\alpha \geq 0$.

(vi) *Duality for the Rényi conditional entropy:* Let $|\psi_{ABC}\rangle$ be a pure state and for $\alpha \geq 1/2$ define β through $1/\alpha + 1/\beta = 2$. Then

$$\tilde{S}_\alpha(A|B)_\psi = -\tilde{S}_\beta(A|C)_\psi.$$

(vii) [19, 3] *Data processing inequality:* If $\alpha \in [1/2, \infty)$ and Λ is a CPTP map, then

$$\tilde{D}_\alpha(\rho \| \sigma) \geq \tilde{D}_\alpha(\Lambda(\rho) \| \Lambda(\sigma)).$$

Furthermore, let $\Phi: B \rightarrow C$ be a CPTP map, and for a bipartite state ρ_{AB} set $\sigma_{AC} := (\text{id}_A \otimes \Phi)(\rho_{AB})$. We then have for $\alpha \geq 1/2$ that

$$\begin{aligned}\tilde{S}_\alpha(A|B)_\rho &\leq \tilde{S}_\alpha(A|C)_\sigma, \\ \tilde{I}_\alpha(A; B)_\rho &\geq \tilde{I}_\alpha(A; C)_\sigma.\end{aligned}$$

2.4 Further properties of Rényi entropic quantities

In this section we derive various properties of the Rényi entropy, the Rényi conditional entropy, and the Rényi mutual information (see Definition 2.4), as well as the Rényi conditional mutual information, defined in (2.14) below. We defer the proofs of the statements concerning Rényi mutual information and Rényi conditional mutual information to Appendix A.

We start with the following subadditivity property for the Rényi entropies [55], for which we give a simplified proof based on the data processing inequality.

Lemma 2.6 (Subadditivity of Rényi entropies [55]). *If $\alpha \geq 0$ and $\rho_{AB} \in \mathcal{D}(\mathcal{H}_{AB})$, then⁵*

$$S_\alpha(A)_\rho - \log |B| \leq S_\alpha(AB)_\rho \leq S_\alpha(A)_\rho + \log |B|.$$

Proof. To prove the upper bound on $S_\alpha(AB)_\rho$, observe first that

$$S_\alpha(AB)_\rho = -\tilde{D}_\alpha(\rho_{AB} \| I_A \otimes \pi_B) + \log |B| \tag{2.7}$$

$$= -D_\alpha(\rho_{AB} \| I_A \otimes \pi_B) + \log |B|. \tag{2.8}$$

Assuming that $\alpha \geq 1/2$ and using (2.7), we then have

$$S_\alpha(AB)_\rho = -\tilde{D}_\alpha(\rho_{AB} \| I_A \otimes \pi_B) + \log |B|$$

⁵Note that both inequalities in Lemma 2.6 can be tightened by replacing the log terms with the 0-Rényi entropy $S_0(B)_\rho = \log \text{rk } \rho_B$. However, throughout this paper we assume Hilbert spaces to be restricted to the support of the corresponding quantum states, so that $S_0(B)_\rho = \log |B|$.

$$\begin{aligned}
&\leq -\tilde{D}_\alpha(\rho_A \| I_A) + \log |B| \\
&= S_\alpha(A)_\rho + \log |B|,
\end{aligned}$$

where the inequality follows from the data processing inequality (Proposition 2.5(vii)). If $\alpha \in [0, 1/2)$, we use relation (2.8) together with the data processing inequality for the α -Rényi relative entropy [34] instead. The lower bound on $S_\alpha(AB)_\rho$ follows from the upper bound by a simple duality argument, using Proposition 2.5(v), as discussed in [55]. \square

We proceed with the following lemma concerning dimension bounds on the Rényi conditional entropy and mutual information, as well as invariance properties with respect to tensor product states.

Lemma 2.7. *Let $\alpha \in [1/2, \infty)$.*

(i) *For an arbitrary tripartite state ρ_{ABC} we have*

$$\tilde{S}_\alpha(A|BC)_\rho + 2 \log |C| \geq \tilde{S}_\alpha(A|B)_\rho, \quad (2.9)$$

$$\tilde{I}_\alpha(A; B)_\rho + 2 \log |C| \geq \tilde{I}_\alpha(A; BC)_\rho. \quad (2.10)$$

(ii) *For states ρ_{AB} and σ_C , we have*

$$\tilde{S}_\alpha(A|BC)_{\rho \otimes \sigma} = \tilde{S}_\alpha(A|B)_\rho, \quad (2.11)$$

$$\tilde{I}_\alpha(A; BC)_{\rho \otimes \sigma} = \tilde{I}_\alpha(A; B)_\rho. \quad (2.12)$$

Proof. We first prove (2.9). By [30, Prop. 8] we have the following bound:

$$\tilde{S}_\alpha(A|BC)_\rho \geq \tilde{S}_\alpha(AC|B)_\rho - \log |C| \quad (2.13)$$

Now consider from duality (Proposition 2.5(vi)) that

$$\tilde{S}_\alpha(AC|B)_\rho = -\tilde{S}_\beta(AC|D)_\rho,$$

where D is a purifying system and β is such that $1/\alpha + 1/\beta = 2$. By the same reasoning, we find that

$$\tilde{S}_\beta(A|CD)_\rho \geq \tilde{S}_\beta(AC|D)_\rho - \log |C|.$$

But from duality this is the same as

$$\tilde{S}_\alpha(A|B)_\rho - \log |C| \leq \tilde{S}_\alpha(AC|B)_\rho.$$

Substituting this in (2.13) then yields the claim.

We continue with the proof of (2.11). From the data processing inequality (Proposition 2.5(vii)), we know that

$$\tilde{S}_\alpha(A|BC)_{\rho \otimes \sigma} \leq \tilde{S}_\alpha(A|B)_\rho.$$

On the other hand, consider that

$$\begin{aligned}
-\tilde{S}_\alpha(A|BC)_{\rho \otimes \sigma} &= \min_{\tau_{BC}} \tilde{D}_\alpha(\rho_{AB} \otimes \sigma_C \| I_A \otimes \tau_{BC}) \\
&\leq \tilde{D}_\alpha(\rho_{AB} \otimes \sigma_C \| I_A \otimes \theta_B \otimes \sigma_C) \\
&= \tilde{D}_\alpha(\rho_{AB} \| I_A \theta_B).
\end{aligned}$$

Since the inequality holds for all θ_B , we get that

$$-\tilde{S}_\alpha(A|BC)_{\rho \otimes \sigma} \leq -\tilde{S}_\alpha(A|B)_\rho,$$

which yields (2.11).

The corresponding relations (2.10) and (2.12) for the Rényi mutual information are proved in Appendix A. \square

The following proposition is crucial for our proofs. It bounds the difference of Rényi entropic quantities of two quantum states in terms of their fidelity. Note that the inequality (2.15) for the Rényi entropies originally appeared in [55]. The last assertion concerns the Rényi conditional mutual information $\tilde{I}_\alpha(A; B|C)_\rho$, defined in [10] for a tripartite state ρ_{ABC} and $\alpha > 0$ as

$$\tilde{I}_\alpha(A; B|C)_\rho = \frac{2\alpha}{\alpha - 1} \log \left\| \rho_{ABC}^{1/2} \rho_{AC}^{(1-\alpha)/2\alpha} \rho_C^{(\alpha-1)/2\alpha} \rho_{BC}^{(1-\alpha)/2\alpha} \right\|_{2\alpha}. \quad (2.14)$$

Note that this quantity does not feature in our proofs. However, the corresponding fidelity bound in (2.18) might be of independent interest, and we include it for the sake of completeness.

Proposition 2.8. *Let $\rho_{AB}, \sigma_{AB} \in \mathcal{D}(\mathcal{H}_{AB})$, and for $\alpha \in (1/2, 1)$ define $\beta \equiv \beta(\alpha) := \alpha/(2\alpha - 1)$. Then the following inequalities hold:*

$$S_\alpha(A)_\rho - S_\beta(A)_\sigma \geq \frac{2\alpha}{1-\alpha} \log F(\rho_A, \sigma_A), \quad (2.15)$$

$$\tilde{S}_\alpha(A|B)_\rho - \tilde{S}_\beta(A|B)_\sigma \geq \frac{2\alpha}{1-\alpha} \log F(\rho_{AB}, \sigma_{AB}). \quad (2.16)$$

Assuming that $\rho_A = \sigma_A$ holds, we also have

$$\tilde{I}_\beta(A; B)_\rho - \tilde{I}_\alpha(A; B)_\sigma \geq \frac{2\alpha}{1-\alpha} \log F(\rho_{AB}, \sigma_{AB}). \quad (2.17)$$

Let ρ_{ABC} and σ_{ABC} be tripartite states satisfying $\rho_{AC} = \sigma_{AC}$, $\rho_{BC} = \sigma_{BC}$, and $\rho_C = \sigma_C$. Furthermore, assume that ρ_{BC} has full support. Then

$$\tilde{I}_\beta(A; B|C)_\rho - \tilde{I}_\alpha(A; B|C)_\sigma \geq \frac{2\alpha}{1-\alpha} \log F(\rho_{ABC}, \sigma_{ABC}). \quad (2.18)$$

Proof. We first observe that (2.15) follows from (2.16) by setting $B = C$. To prove (2.16), let τ_B be an arbitrary density operator, and let $\varepsilon \in (0, 1)$. Furthermore, let

$$\tau(\varepsilon)_B := (1 - \varepsilon)\tau_B + \varepsilon\rho_B,$$

and recall that

$$\tilde{D}_\alpha(\omega||\theta) \leq \tilde{D}_\alpha(\omega||\theta') \quad (2.19)$$

holds for all $\alpha \in [1/2, \infty]$ and for $\theta \geq \theta' \geq 0$ [30, Prop. 4]. Observe also that for $c > 0$ we have

$$\tilde{D}_\alpha(\omega||c\theta) = \tilde{D}_\alpha(\omega||\theta) - \log c. \quad (2.20)$$

Consider then the following chain of inequalities:

$$-\tilde{D}_\alpha(\rho_{AB}||I_A \otimes \tau(\varepsilon)_B) + \tilde{D}_\beta(\sigma_{AB}||I_A \otimes \tau_B) - \log(1 - \varepsilon)$$

$$\begin{aligned}
&= -\tilde{D}_\alpha(\rho_{AB} \| I_A \otimes \tau(\varepsilon)_B) + \tilde{D}_\beta(\sigma_{AB} \| I_A \otimes (1 - \varepsilon) \tau_B) \\
&\geq -\tilde{D}_\alpha(\rho_{AB} \| I_A \otimes \tau(\varepsilon)_B) + \tilde{D}_\beta(\sigma_{AB} \| I_A \otimes \tau(\varepsilon)_B) \\
&= \frac{2\alpha}{1 - \alpha} \log \left\| \rho_{AB}^{1/2} \tau(\varepsilon)_B^{(1-\alpha)/2\alpha} \right\|_{2\alpha} + \frac{2\beta}{\beta - 1} \log \left\| \tau(\varepsilon)_B^{(1-\beta)/2\beta} \sigma_{AB}^{1/2} \right\|_{2\beta} \\
&= \frac{2\alpha}{1 - \alpha} \log \left[\left\| \rho_{AB}^{1/2} \tau(\varepsilon)_B^{(1-\alpha)/2\alpha} \right\|_{2\alpha} \left\| \tau(\varepsilon)_B^{(1-\beta)/2\beta} \sigma_{AB}^{1/2} \right\|_{2\beta} \right] \\
&\geq \frac{2\alpha}{1 - \alpha} \log \left\| \rho_{AB}^{1/2} \tau(\varepsilon)_B^{(1-\alpha)/2\alpha} \tau(\varepsilon)_B^{(1-\beta)/2\beta} \sigma_{AB}^{1/2} \right\|_1 \\
&= \frac{2\alpha}{1 - \alpha} \log \left\| \rho_{AB}^{1/2} \sigma_{AB}^{1/2} \right\|_1 \\
&= \frac{2\alpha}{1 - \alpha} \log F(\rho_{AB}, \sigma_{AB}). \tag{2.21}
\end{aligned}$$

The first equality is an application of (2.20). The first inequality is a consequence of (2.19) and the fact that

$$(1 - \varepsilon) \tau_B \leq (1 - \varepsilon) \tau_B + \varepsilon \pi_B = \tau(\varepsilon)_B.$$

The second and third equality follow from the definition of the sandwiched Rényi divergence (see Definition 2.4) and the relation

$$\frac{\beta}{\beta - 1} = \frac{\alpha}{1 - \alpha}. \tag{2.22}$$

The second inequality is an application of Hölder's inequality (2.2) (note that $1/2\alpha + 1/2\beta = 1$). The second-to-last equality follows because $\tau(\varepsilon)_B$ is a full rank operator for $\varepsilon \in (0, 1)$, so that $\tau(\varepsilon)_B^{(1-\alpha)/2\alpha} \tau(\varepsilon)_B^{(1-\beta)/2\beta} = I_B$.

Since (2.21) holds for an arbitrary density operator τ_B , we may choose τ_B to be the optimizing state for $\tilde{S}_\beta(A|B)_\sigma$. We can then continue from (2.21) as

$$\begin{aligned}
\frac{2\alpha}{1 - \alpha} \log F(\rho_{AB}, \sigma_{AB}) &\leq -\tilde{D}_\alpha(\rho_{AB} \| I_A \otimes \tau(\varepsilon)_B) + \tilde{D}_\beta(\sigma_{AB} \| I_A \otimes \tau_B) - \log(1 - \varepsilon) \\
&\leq \max_{\omega_B \in \mathcal{D}(\mathcal{H}_B)} \left\{ -\tilde{D}_\alpha(\rho_{AB} \| I_A \otimes \omega_B) \right\} + \tilde{D}_\beta(\sigma_{AB} \| I_A \otimes \tau_B) - \log(1 - \varepsilon) \\
&= \tilde{S}_\alpha(A|B)_\rho - \tilde{S}_\beta(A|B)_\sigma - \log(1 - \varepsilon).
\end{aligned}$$

We have shown that this relation holds for all $\varepsilon \in (0, 1)$, and so taking the limit $\varepsilon \searrow 0$ yields the claim.

The bounds (2.17) and (2.18) are proved in Appendix A. \square

The next result concerns the Rényi conditional entropy of c-q states. Note that a special case of (2.23) for the Rényi entropies (i.e. where system B is trivial) appeared in [41].

Proposition 2.9. *Let $\rho_{ABX} = \sum_{x \in \mathcal{X}} p_x \rho_{AB}^x \otimes |x\rangle\langle x|_X$ be a c-q state. Then the following properties hold for all $\alpha > 0$:*

(i) *Monotonicity under discarding classical information:*

$$\tilde{S}_\alpha(AX|B)_\rho \geq \tilde{S}_\alpha(A|B)_\rho. \tag{2.23}$$

(ii) *Dimension bound:*

$$\tilde{S}_\alpha(A|BX)_\rho + \log |X| \geq \tilde{S}_\alpha(A|B)_\rho, \quad (2.24)$$

$$\tilde{I}_\alpha(A; BX)_\rho \leq \log |X| + \tilde{I}_\alpha(A; B)_\rho. \quad (2.25)$$

Proof. To prove (2.23), let τ_B be the optimizing state for $\tilde{S}_\alpha(A|B)_\rho$, and assume $\alpha \neq 1$. We then have

$$\begin{aligned} \tilde{S}_\alpha(A|BX)_\rho &\geq \frac{1}{1-\alpha} \log \text{Tr} \left[\left(\tau_B^{(1-\alpha)/2\alpha} \rho_{ABX} \tau_B^{(1-\alpha)/2\alpha} \right)^\alpha \right] \\ &= \frac{1}{1-\alpha} \log \left\{ \sum_{x \in \mathcal{X}} \text{Tr} \left[\left(\tau_B^{(1-\alpha)/2\alpha} p_x \rho_{AB}^x \tau_B^{(1-\alpha)/2\alpha} \right)^\alpha \right] \right\} \\ &= \frac{1}{1-\alpha} \log \left\{ \sum_{x \in \mathcal{X}} \left\| \tau_B^{(1-\alpha)/2\alpha} p_x \rho_{AB}^x \tau_B^{(1-\alpha)/2\alpha} \right\|_\alpha^\alpha \right\} \\ &\geq \frac{1}{1-\alpha} \log \left\{ \left\| \tau_B^{(1-\alpha)/2\alpha} \sum_{x \in \mathcal{X}} p_x \rho_{AB}^x \tau_B^{(1-\alpha)/2\alpha} \right\|_\alpha^\alpha \right\} \\ &= \frac{1}{1-\alpha} \log \left\{ \left\| \tau_B^{(1-\alpha)/2\alpha} \rho_{AB} \tau_B^{(1-\alpha)/2\alpha} \right\|_\alpha^\alpha \right\} \\ &= \tilde{S}_\alpha(A|B)_\rho, \end{aligned}$$

where the inequality follows from McCarthy's inequalities, using (2.3) for $\alpha < 1$ and (2.4) for $\alpha > 1$. For $\alpha = 1$, the claim follows easily from definition (2.1) of the quantum relative entropy.

To prove (2.24), observe that we have

$$\tilde{S}_\alpha(A|BX)_\rho + \log |X| \geq \tilde{S}_\alpha(A|B)_\rho \geq \tilde{S}_\alpha(A|B)_\rho,$$

where the first inequality is [30, Prop. 8], and the second inequality is (2.23).

Finally, (2.25) follows from (2.24) and the reasoning in Appendix A.1. \square

3 State redistribution

3.1 The protocol

Consider a tripartite state ρ_{ABC} shared between Alice and Bob, with the systems A and C being with Alice and the system B being with Bob. Let ψ_{ABCR} denote a purification of ρ_{ABC} , with R being an inaccessible, purifying reference system. Furthermore, Alice and Bob share entanglement in the form of an MES $\Phi_{T_A T_B}^k$ of Schmidt rank k , with the systems T_A and T_B being with Alice and Bob, respectively. The goal of the state redistribution protocol is to transfer the system A from Alice to Bob, while preserving its correlations with the other systems. In the process, the shared entanglement is transformed to an MES $\Phi_{T'_A T'_B}^m$ of Schmidt rank m , where T'_A and T'_B are with Alice and Bob, respectively. The initial state and the target state are shown in Figure 2.

In achieving this goal, Alice and Bob are allowed to use local encoding and decoding operations on the systems in their possession. In addition, Alice is allowed to send qubits to Bob (through a noiseless quantum channel). A general state redistribution protocol (ρ, Λ) with $\Lambda \equiv \mathcal{D} \circ \mathcal{E}$ and $\rho = \rho_{ABC}$ consists of the following steps (cf. Figure 3):

1. Alice applies an encoding CPTP map $\mathcal{E}: ACT_A \rightarrow C'T'_A Q$ and sends the system Q to Bob through the noiseless quantum channel.

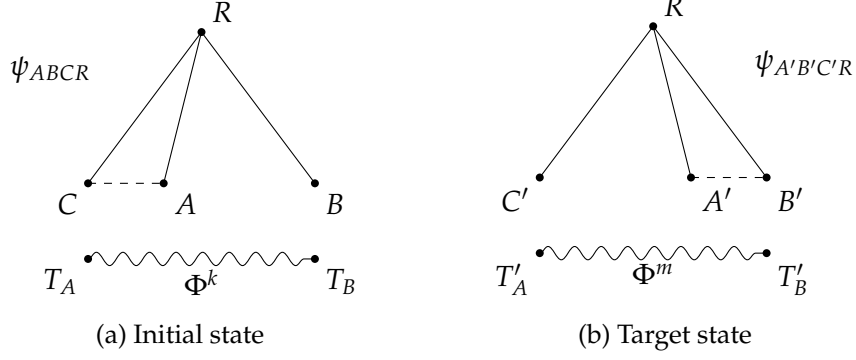


Figure 2: State redistribution protocol that transfers Alice’s system A to Bob. Starting with the initial state $\psi \otimes \Phi^k$ depicted in (a), the protocol outputs a state that is close in fidelity to the target state $\psi \otimes \Phi^m$ depicted in (b).

2. Upon receiving the system Q , Bob applies a decoding CPTP map $\mathcal{D}: QBT_B \rightarrow T'_B A' B'$, where $T'_B \cong T'_A$ and $A' \cong A$.

The initial state shared between Alice, Bob, and the reference is $\Phi_{T_A T_B}^k \otimes \psi_{ABCR}$, the state after Alice’s encoding operation is

$$\omega \equiv \omega_{T'_A T'_B C' Q B R} := (\mathcal{E} \otimes \text{id}_R)(\Phi^k \otimes \psi), \quad (3.1)$$

and the final state of the protocol (ρ, Λ) is given by

$$\sigma \equiv \sigma_{T'_A T'_B A' B' C' R} := (\Lambda \otimes \text{id}_R)(\Phi^k \otimes \psi) = (\mathcal{D} \circ \mathcal{E} \otimes \text{id}_R)(\Phi^k \otimes \psi). \quad (3.2)$$

The aim is to obtain a state σ that is close to the target state $\Phi_{T'_A T'_B}^m \otimes \psi_{A' B' C' R}$, where $\psi_{A' B' C' R} = \psi_{ABCR}$. The figure of merit of the protocol is the fidelity $F(\sigma, \Phi^m \otimes \psi)$. The number of qubits that Alice sends to Bob, is given by $\log |Q|$, whereas the number of ebits consumed in the protocol is given by $\log k - \log m = \log |T_A| - \log |T'_A|$. If $k < m$ then ebits are *gained* in the protocol.

We consider state redistribution (and all other tasks studied in this paper) in the asymptotic, memoryless setting, where Alice and Bob start with multiple (say, n) copies of the initial resource, and the strong converse property is established in the limit $n \rightarrow \infty$. In this case, Alice and Bob initially share n identical copies of the state ρ_{ABC} with purification ψ_{ABCR} , i.e. they share the state $\rho_{ABC}^{\otimes n}$ with purification $\psi_{ABCR}^{\otimes n}$. Moreover, they share an MES $\Phi_{T'_A T'_B}^{k_n}$ of Schmidt rank k_n . We then consider state redistribution protocols $(\rho^{\otimes n}, \Lambda_n)$ with the figure of merit

$$F_n := F(\sigma_n, \Phi^{m_n} \otimes \psi^{\otimes n}), \quad (3.3)$$

where $\Phi_{T'_A T'_B}^{m_n}$ is an MES of Schmidt rank m_n , and $\sigma_n := (\Lambda_n \otimes \text{id}_R)(\Phi^{k_n} \otimes \psi^{\otimes n})$, where $\Lambda_n: A^n C^n T'_A^n \otimes B^n T'_B^n \rightarrow C'^n T'^n \otimes T'_B'^n A'^n B'^n$ with Q^n being sent from Alice to Bob. The two operational quantities of interest are as follows:

1. The quantum communication cost of the protocol $(\rho^{\otimes n}, \Lambda_n)$, given by

$$q(\rho^{\otimes n}, \Lambda_n) := \frac{1}{n} \log |Q^n|. \quad (3.4)$$

2. The entanglement cost of the protocol $(\rho^{\otimes n}, \Lambda_n)$, given by

$$e(\rho^{\otimes n}, \Lambda_n) := \frac{1}{n} (\log k_n - \log m_n) = \frac{1}{n} (\log |T'_A^n| - \log |T'_A'^n|). \quad (3.5)$$

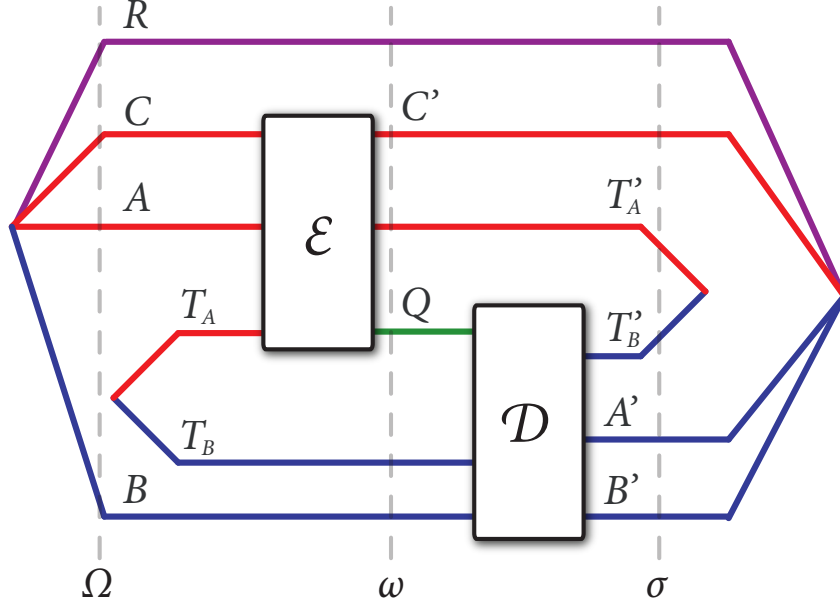


Figure 3: State redistribution protocol (see Section 3.1 for a detailed explanation).

A pair $(e, q) \in \mathbb{R}^2$ with $q \geq 0$, is said to be an achievable rate pair for state redistribution of a state ρ_{ABC} , if there exists a sequence of protocols $\{(\rho^{\otimes n}, \Lambda_n)\}_{n \in \mathbb{N}}$, satisfying

$$\limsup_{n \rightarrow \infty} q(\rho^{\otimes n}, \Lambda_n) = q, \quad \limsup_{n \rightarrow \infty} e(\rho^{\otimes n}, \Lambda_n) = e, \quad \liminf_{n \rightarrow \infty} F_n = 1.$$

Luo and Devetak [27] and Yard and Devetak [54] (see also Devetak and Yard [17]) proved that a pair (e, q) is an achievable rate pair for state redistribution of a state ρ_{ABC} if and only if it lies in the region (cf. Figure 1) defined by

$$q \geq \frac{1}{2}I(A; R|B)_\rho, \quad q + e \geq S(A|B)_\rho. \quad (3.6)$$

Recently, a strong converse theorem for the quantum communication cost was proved by Berta *et al.* [8], using the smooth entropy framework (cf. [38, 43, 12, 44] and references therein). This theorem, however, did not prove the strong converse property for the entire boundary of the achievable rate region given by (3.6). We fill this gap with Theorem 3.2, as well as provide an alternative proof of the strong converse theorem of [8]. As mentioned earlier, Berta *et al.* have now also extended their proof to the entire boundary of the achievable rate region [9].

3.2 Strong converse theorem

Lemma 3.1. *Let $\rho \equiv \rho_{ABC}$ be a tripartite state with purification $|\psi_{ABCR}\rangle$, and let (ρ, Λ) be a state redistribution protocol where $\Lambda \equiv \mathcal{D} \circ \mathcal{E}$ with $\mathcal{E}: ACT_A \rightarrow C'T'_AQ$ and $\mathcal{D}: QBT_B \rightarrow T'_B A'B'$ as defined in Section 3.1. Furthermore, set*

$$F := F\left(\Phi_{T'_A T'_B}^m \otimes \psi_{A'B'C'R}, (\mathcal{D} \circ \mathcal{E} \otimes \text{id}_R)\left(\Phi_{T_A T_B}^k \otimes \psi_{ABCR}\right)\right).$$

Then we have the following bounds on F for $\alpha \in (1/2, 1)$ and $\beta \equiv \beta(\alpha) = \alpha/(2\alpha - 1)$:

$$\log F \leq \frac{1-\alpha}{2\alpha} (\log |Q| + \log |T_A| - \log |T'_A| - S_\beta(AB)_\rho + S_\alpha(B)_\rho), \quad (3.7)$$

$$\log F \leq \frac{1-\alpha}{2\alpha} \left(2\log |Q| - \tilde{S}_\beta(R|B)_\rho + \tilde{S}_\alpha(R|AB)_\rho \right). \quad (3.8)$$

We also have the following alternative bound to (3.8):

$$\log F \leq \frac{1-\alpha}{2\alpha} \left(2\log |Q| - \tilde{I}_\alpha(R; AB)_\rho + \tilde{I}_\beta(R; B)_\rho \right). \quad (3.9)$$

Proof. We first prove (3.7). Denote by $U_{\mathcal{E}}: \mathcal{H}_{ACT_A} \rightarrow \mathcal{H}_{C'T'_A Q E_1}$ and $U_{\mathcal{D}}: \mathcal{H}_{QBT_B} \rightarrow \mathcal{H}_{T'_B A' B' E_2}$ the Stinespring isometries of the maps \mathcal{E} and \mathcal{D} respectively, and define the pure states

$$|\omega_{C'T'_A Q T_B B R E_1}\rangle := U_{\mathcal{E}} \left(|\Phi_{T_A T_B}^k\rangle \otimes |\psi_{ABCR}\rangle \right), \quad (3.10)$$

$$|\sigma_{T'_A T'_B A' B' C' R E_1 E_2}\rangle := U_{\mathcal{D}} |\omega_{C'T'_A Q T_B B R E_1}\rangle, \quad (3.11)$$

that purify the mixed states ω and σ defined in (3.1) and (3.2), respectively. We then have

$$S_\alpha(QT_B B)_\omega \leq \log |Q| + \log |T_A| + S_\alpha(B)_\rho, \quad (3.12)$$

where we used the subadditivity of the Rényi entropies (Lemma 2.6) twice, as well as the fact that T_A is the same size as T_B . For the fidelity $F := F(\Phi_{T'_A T'_B}^m \otimes \psi_{A'B'C'R}, \sigma_{T'_A T'_B A' B' C' R})$, we know by Uhlmann's theorem that there exists a pure state $\phi_{E_1 E_2}$ such that

$$\begin{aligned} F &= F \left(\Phi_{T'_A T'_B}^m \otimes \psi_{A'B'C'R} \otimes \phi_{E_1 E_2}, \sigma_{T'_A T'_B A' B' C' R E_1 E_2} \right) \\ &\leq F(\pi_{T'_B}^m \otimes \rho_{A'B'} \otimes \phi_{E_2}, \sigma_{T'_B A' B' E_2}), \end{aligned}$$

where $|\sigma_{T'_A T'_B A' B' C' R E_1 E_2}\rangle$ is the pure state defined in (3.11). The inequality follows from the monotonicity of the fidelity under partial trace. Hence, by eq. (2.15) of Proposition 2.8 we obtain the following bound, setting $\beta = \alpha/(2\alpha - 1)$:

$$\begin{aligned} S_\alpha(QT_B B)_\omega &= S_\alpha(T'_B A' B' E_2)_\sigma \\ &\geq S_\beta(T'_B A' B' E_2)_{\pi^m \otimes \rho \otimes \phi} + \frac{2\alpha}{1-\alpha} \log F \\ &\geq \log |T'_A| + S_\beta(A'B')_\rho + \frac{2\alpha}{1-\alpha} \log F, \end{aligned} \quad (3.13)$$

where we used the invariance of the Rényi entropies under the isometry $U_{\mathcal{D}}$ (Proposition 2.5) in the first equality, eq. (2.15) of Proposition 2.8 in the first inequality, and additivity and positivity of the Rényi entropies (Proposition 2.5(iii) and (ii)) in the second inequality. Combining (3.12) and (3.13) then yields

$$\log |Q| + \log |T_A| + S_\alpha(B)_\rho \geq \log |T'_A| + S_\beta(AB)_\rho + \frac{2\alpha}{1-\alpha} \log F,$$

which is equivalent to (3.7).

To prove (3.8), we first note that from eq. (2.16) of Proposition 2.8 we have the inequality

$$\frac{2\alpha}{1-\alpha} \log F(\rho_{ABR}, \sigma_{A'B'R}) \leq \tilde{S}_\alpha(R|AB)_\rho - \tilde{S}_\beta(R|A'B')_\sigma. \quad (3.14)$$

We bound the second term on the right-hand side of (3.14) as follows:

$$-\tilde{S}_\beta(R|A'B')_\sigma \leq -\tilde{S}_\beta(R|QBT_B)_\omega$$

$$\begin{aligned}
&\leq -\tilde{S}_\beta(R|BT_B)_\omega + 2\log|Q| \\
&= -\tilde{S}_\beta(R|B)_\omega + 2\log|Q| \\
&= -\tilde{S}_\beta(R|B)_\rho + 2\log|Q|,
\end{aligned} \tag{3.15}$$

where we used data processing (Proposition 2.5(vii)) in the first inequality, and eq. (2.9) and (2.11) of Lemma 2.7 in the second inequality and the first equality, respectively. Substituting (3.15) in (3.14) now yields (3.8).

The bound (3.9) follows from similar arguments as those used for the proof of (3.8), relying on eq. (2.10) and (2.12) of Lemma 2.7 and eq. (2.17) of Proposition 2.8 instead. We therefore omit an explicit proof. \square

Lemma 3.1 immediately implies the following strong converse theorem:

Theorem 3.2 (Strong converse for state redistribution). *Let $\rho \equiv \rho_{ABC}$ be a tripartite state and $\{(\rho^{\otimes n}, \Lambda_n)\}_{n \in \mathbb{N}}$ be a sequence of state redistribution protocols as described in Section 3.1, with figure of merit F_n as defined in (3.3). Then for all $n \in \mathbb{N}$ we have the following bounds on F_n for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:*

$$F_n \leq \exp \left\{ -n\kappa(\alpha) \left[S_\beta(AB)_\rho - S_\alpha(B)_\rho - (q + e) \right] \right\}, \tag{3.16}$$

$$F_n \leq \exp \left\{ -n\kappa(\alpha) \left[\tilde{S}_\beta(R|B)_\rho - \tilde{S}_\alpha(R|AB)_\rho - 2q \right] \right\}, \tag{3.17}$$

where $\kappa(\alpha) = (1 - \alpha)/(2\alpha)$, and $q \equiv q(\rho^{\otimes n}, \Lambda_n)$ and $e \equiv e(\rho^{\otimes n}, \Lambda_n)$ are the quantum communication cost and entanglement cost defined in (3.4) and (3.5), respectively. As an alternative to (3.17), we also obtain the bound

$$F_n \leq \exp \left\{ -n\kappa(\alpha) \left[\tilde{I}_\alpha(R; AB)_\rho - \tilde{I}_\beta(R; B)_\rho - 2q \right] \right\}. \tag{3.18}$$

We have

$$S_\beta(AB)_\rho - S_\alpha(B)_\rho \xrightarrow{\alpha \rightarrow 1} S(A|B)_\rho, \tag{3.19a}$$

$$\tilde{S}_\beta(R|B)_\rho - \tilde{S}_\alpha(R|AB)_\rho \xrightarrow{\alpha \rightarrow 1} I(A; R|B)_\rho, \tag{3.19b}$$

where (3.19b) follows from a similar argument as in [11, Lem. 10]. Moreover, the right-hand sides of (3.19) determine the boundary of the region of achievable rate pairs (e, q) given by (3.6). Hence, we obtain the following strong converse theorem along the same lines as at the end of Section 1.1: If $q + e < S(A|B)_\rho$ or $q < \frac{1}{2}I(A; R|B)_\rho$, there is a constant $K > 0$ such that

$$F_n \leq \exp(-nK).$$

For the remainder of the paper, we will skip this last step, and merely state strong converse theorems in the form of Theorem 3.2.

3.3 Rényi generalizations of the conditional mutual information

We can regard the expressions $\tilde{S}_\beta(R|B)_\rho - \tilde{S}_\alpha(R|AB)_\rho$ and $\tilde{I}_\alpha(R; AB)_\rho - \tilde{I}_\beta(R; B)_\rho$ appearing in the bounds on the fidelity in Lemma 3.1 and Theorem 3.2 as Rényi generalizations of the conditional mutual information $I(A; R|B)_\rho$ (see [10] for a detailed discussion of this concept). More generally, for a tripartite state ρ_{ABC} and Rényi parameters $\alpha \geq 1/2$ and $\beta \equiv \beta(\alpha) = \alpha/(2\alpha - 1)$, we define

$$\tilde{I}_\alpha^{(1)}(A; B|C)_\rho := \tilde{S}_\alpha(A|C)_\rho - \tilde{S}_\beta(A|BC)_\rho, \tag{3.20a}$$

$$\tilde{I}_\alpha^{(2)}(A; B|C)_\rho := \tilde{I}_\alpha(A; BC)_\rho - \tilde{I}_\beta(A; C)_\rho. \quad (3.20b)$$

These quantities satisfy the following properties:

Proposition 3.3. *Let ρ_{ABC} be a tripartite state and $\alpha \geq 1/2$. The quantities $\tilde{I}_\alpha^{(i)}(A; B|C)_\rho$, defined by (3.20) for $i = 1, 2$, satisfy:*

(i) *Rényi generalization of conditional mutual information:*

$$\lim_{\alpha \rightarrow 1} \tilde{I}_\alpha^{(i)}(A; B|C)_\rho = I(A; B|C)_\rho.$$

(ii) *Monotonicity in α : For $1/2 \leq \alpha \leq \alpha'$, we have*

$$\tilde{I}_\alpha^{(1)}(A; B|C)_\rho \geq \tilde{I}_{\alpha'}^{(1)}(A; B|C)_\rho \quad \tilde{I}_\alpha^{(2)}(A; B|C)_\rho \leq \tilde{I}_{\alpha'}^{(2)}(A; B|C)_\rho.$$

(iii) *Data processing inequality on the B system: Let $\Lambda: B \rightarrow B'$ be a CPTP map and define $\sigma_{AB'C} = (\text{id}_{AC} \otimes \Lambda)(\rho_{ABC})$, then*

$$\tilde{I}_\alpha^{(i)}(A; B|C)_\rho \geq \tilde{I}_\alpha^{(i)}(A; B'|C)_\sigma.$$

(iv) *Duality: Let ρ_{ABCD} be a purification of ρ_{ABC} , then*

$$\tilde{I}_\alpha^{(1)}(A; B|C)_\rho = \tilde{I}_\alpha^{(1)}(A; B|D)_\rho.$$

Proof. Property (i) follows from [11, Lem. 8]. For (ii), note that $\beta(\alpha) = \alpha/(2\alpha - 1)$ is a decreasing function. The assertion then follows from the monotonicity in α of the sandwiched Rényi divergence, Proposition 2.5(i).⁶ Property (iii) is straightforward, and (iv) is obtained by employing duality (Proposition 2.5(vi)). \square

3.4 State redistribution with feedback

In this section, we consider state redistribution with feedback [8], where the state redistribution protocol consists of M rounds of forward and backward quantum communication between Alice and Bob. The initial state of the protocol is again the pure state $\psi_{ABCR} \otimes \Phi_{T_A T_B}^k$, where systems A and C are with Alice, B is with Bob, R is an inaccessible reference system, and $\Phi_{T_A T_B}^k$ is an MES of Schmidt rank k shared between Alice (T_A) and Bob (T_B). As before, the goal is for Alice to transfer the A system to Bob, while preserving its correlations with the other systems.

The main difference with single-round state redistribution as described in Section 3.1 is that now backward quantum communication from Bob to Alice is possible. Furthermore, we allow for M rounds of communication, in the following way (cf. Figure 4): Alice first applies an encoding operation $\mathcal{E}_1: ACT_A \rightarrow Q_1 A_1$ to the initial state and sends Q_1 to Bob, who applies a decoding operation $\mathcal{D}_1: Q_1 B T_B \rightarrow Q'_1 B_1$. The system Q'_1 is the quantum communication register that he sends back to Alice. She then applies the encoding $\mathcal{E}_2: Q'_1 A_1 \rightarrow Q_2 A_2$ and sends Q_2 to Bob, who applies the decoding $\mathcal{D}_2: Q_2 B_1 \rightarrow Q'_2 B_2$ and sends Q'_2 back, and so forth. In the i -th round, we denote by ω^i and σ^i the states shared between Alice, Bob, and the reference, after applying the encoding \mathcal{E}_i and decoding \mathcal{D}_i , respectively. In the final round, Alice applies the encoding $\mathcal{E}_M: Q'_{M-1} A_{M-1} \rightarrow Q_M C' T'_A$ and sends Q_M to Bob, who applies the decoding $\mathcal{D}_M: Q_M B_{M-1} \rightarrow A' B' T'_B$. The protocol succeeds if the final state is close in fidelity to the pure

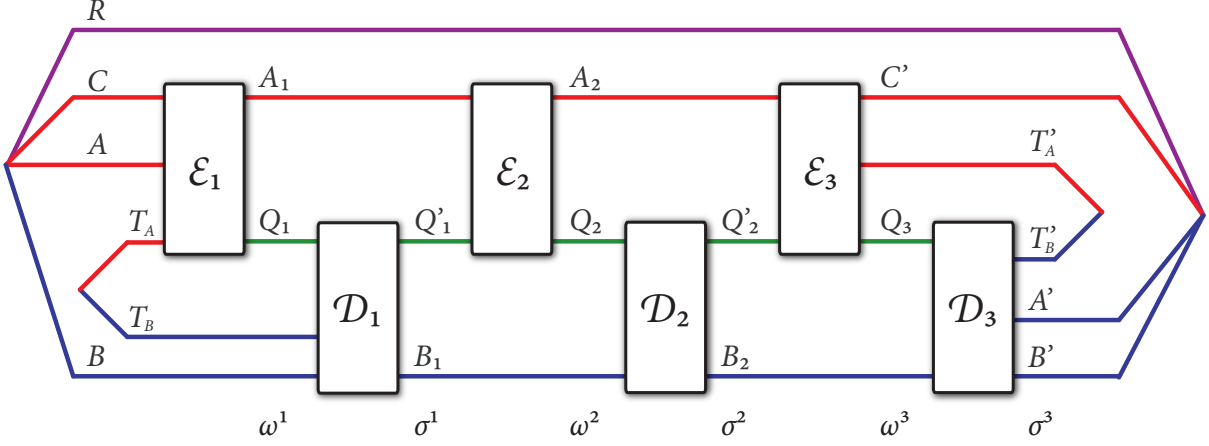


Figure 4: State redistribution protocol with feedback for $M = 3$ (see Section 3.4 for a detailed description).

target state $\psi_{A'B'C'R} \otimes \Phi_{T'_A T'_B}^m$, where $\psi_{A'B'C'R} = \psi_{ABCR}$ and $\Phi_{T'_A T'_B}^m$ is an MES of Schmidt rank m (for some $m \in \mathbb{N}$) shared between Alice and Bob.

For a protocol acting on a many-copy initial state $\psi_{ABCR}^{\otimes n} \otimes \Phi_{T'_A T'_B}^{k_n}$, we define the entanglement cost e , the forward quantum communication q_{\rightarrow} , and the total quantum communication q_{\leftrightarrow} (equal to forward plus backward communication) as

$$e := \frac{1}{n} (\log |T'_A|^n - \log |T''_A|^n), \quad (3.21)$$

$$q_{\rightarrow} := \frac{1}{n} \sum_{i=1}^M \log |Q_i^n|, \quad (3.22)$$

$$q_{\leftrightarrow} := q_{\rightarrow} + \frac{1}{n} \sum_{i=1}^{M-1} \log |Q'_i|^n. \quad (3.23)$$

Using the smooth entropy framework, the authors in [8] proved that the achievable region for e , q_{\rightarrow} , and q_{\leftrightarrow} is determined by the following conditions:

$$q_{\rightarrow} \geq \frac{1}{2} I(A; R|B)_{\psi}, \quad q_{\leftrightarrow} + e \geq S(A|B)_{\psi}. \quad (3.24)$$

We first note that both conditions are independent of M , the number of rounds. Let us also compare (3.24) to the conditions for single-round state redistribution in (3.6). The achievable region for forward quantum communication coincides with that of (3.6), as shown in [8]. However, for state redistribution with feedback the condition for the entanglement cost involves the *total* quantum communication between the two parties. This can be understood to arise from the fact that quantum communication from Bob to Alice can introduce additional entanglement between them. In comparison to a single-round state redistribution protocol without feedback, the optimal rate of the overall entanglement cost e is *lowered* by the amount of backward quantum communication, since we have $q_{\leftrightarrow} \geq q_{\rightarrow}$.

Using the Rényi entropy method, we derive a strong converse theorem for state redistribution with feedback, Theorem 3.4 below. It follows from Lemma B.1, which we state and prove in Appendix B. The strong converse for state redistribution with feedback originally appeared in [8]. However, in contrast to [8] our proof method yields a Rényi entropic quantity as an explicit exponent in the strong converse bound.

⁶Note that the quantity $S_{\beta}(AB)_{\rho} - S_{\alpha}(B)_{\rho}$ that appears in Theorem 3.2 is monotonic in α for the same reason.

Theorem 3.4 (Strong converse for state redistribution with feedback). *Let Λ_n denote a state redistribution protocol with feedback for a pure state $\psi_{ABCR}^{\otimes n}$, as described above. Setting*

$$F_n := F \left(\psi_{A'B'C'R}^{\otimes n} \otimes \Phi_{T_A' T_B'}^{m_n}, (\Lambda_n \otimes \text{id}_{R^n}) \left(\psi_{ABCR}^{\otimes n} \otimes \Phi_{T_A T_B}^{k_n} \right) \right),$$

we then have the following bounds on F_n for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:

$$F_n \leq \exp \left\{ -n\kappa(\alpha) \left[S_\beta(AB)_\psi - S_\alpha(B)_\psi - (q_{\leftrightarrow} + e) \right] \right\}, \quad (3.25)$$

$$F_n \leq \exp \left\{ -n\kappa(\alpha) \left[\tilde{S}_\beta(R|B)_\psi - \tilde{S}_\alpha(R|AB)_\psi - 2q_{\rightarrow} \right] \right\}, \quad (3.26)$$

$$F_n \leq \exp \left\{ -n\kappa(\alpha) \left[\tilde{I}_\alpha(R; AB)_\psi - \tilde{I}_\beta(R; B)_\psi - 2q_{\rightarrow} \right] \right\}, \quad (3.27)$$

where $\kappa(\alpha) = (1 - \alpha)/(2\alpha)$, and e , q_{\rightarrow} , and q_{\leftrightarrow} are the entanglement cost (3.21), forward quantum communication cost (3.22), and total quantum communication cost (3.23), respectively.

3.5 Coherent state merging

Coherent state merging [33] is the task in which Alice wants to transfer the A -part of a bipartite state ρ_{AB} to Bob (who holds the B system), while at the same time generating entanglement between them. We assume that ρ_{AB} is purified by an inaccessible reference system R . To achieve their goal, Alice and Bob are allowed to perform local operations on the systems in their possession as well as noiseless quantum communication. This protocol (also known as ‘Fully Quantum Slepian Wolf’ (FQSW) protocol [1]) is a special case of the state redistribution protocol from Section 3.1 where the system C is absent (or equivalently taken to be a trivial one-dimensional system), and Alice and Bob do not share any entanglement prior to commencing the protocol (hence, the systems T_A and T_B are trivial and entanglement is always *gained* in the course of the protocol).

Let Alice and Bob share n identical copies of the state ρ_{AB} with purification ψ_{ABR} , i.e. the state $\rho_{AB}^{\otimes n}$ with purification $\psi_{ABR}^{\otimes n}$. A general coherent state merging protocol $(\rho^{\otimes n}, \Lambda_n)$ is given by a joint quantum operation $\Lambda_n \equiv \mathcal{D}_n \circ \mathcal{E}_n$ where $\mathcal{E}_n: A^n \rightarrow T_A^{m_n} Q^n$ is Alice’s encoding map, the system Q^n is sent to Bob, and $\mathcal{D}_n: Q^n B^n \rightarrow T_B^{m_n} A^{m_n} B^{m_n}$ is Bob’s decoding map. Here, $A^{m_n} \cong A^n$, $T_B^{m_n} \cong T_A^{m_n}$, and $B^{m_n} \cong B^n$. Denoting the final state of the protocol by $\sigma_n = (\Lambda_n \otimes \text{id}_{R^n})(\psi_{ABR}^{\otimes n})$, the figure of merit is chosen to be the fidelity

$$F_n := F \left(\sigma_n, \psi_{ABR}^{\otimes n} \otimes \Phi_{T_A^{m_n} T_B^{m_n}}^{m_n} \right) \quad (3.28)$$

where the second argument of the fidelity is the target state of the protocol $(\rho^{\otimes n}, \Lambda_n)$.

The *quantum communication cost* $q_{\text{csm}}(\rho^{\otimes n}, \Lambda_n)$ and the *entanglement gain* $e_{\text{csm}}(\rho^{\otimes n}, \Lambda_n)$ are defined in analogy to Section 3.1:

$$q_{\text{csm}}(\rho^{\otimes n}, \Lambda_n) := \frac{1}{n} \log |Q^n|, \quad e_{\text{csm}}(\rho^{\otimes n}, \Lambda_n) := \frac{1}{n} \log |T_A^{m_n}|. \quad (3.29)$$

A pair (e, q) , with $e, q \geq 0$, is said to be an achievable rate pair for coherent state merging of a state ρ_{AB} , if there exists a sequence of protocols $\{(\rho^{\otimes n}, \Lambda_n)\}_{n \in \mathbb{N}}$ such that $\liminf_{n \rightarrow \infty} F_n = 1$ and

$$\limsup_{n \rightarrow \infty} q_{\text{csm}}(\rho^{\otimes n}, \Lambda_n) = q, \quad \liminf_{n \rightarrow \infty} e_{\text{csm}}(\rho^{\otimes n}, \Lambda_n) = e.$$

Coherent state merging was introduced by Abeyesinghe *et al.* [1] and further investigated in [13, 7]. It was proved that a rate pair (e, q) is achievable if and only if e and q satisfy the conditions

$$q \geq \frac{1}{2}I(A; R)_\rho, \quad q - e \geq S(A|B)_\rho.$$

As mentioned above, every coherent state merging protocol can be seen as a special case of a state redistribution protocol where the systems C and T_A are trivial. In this case, $I(A; R|B)_\rho = I(A; R)_\rho$, and Lemma 3.1 reduces to

Lemma 3.5. *Let $\rho \equiv \rho_{AB}$ be a bipartite state with purification $|\psi_{ABR}\rangle$, and let (ρ, Λ) be a coherent state merging protocol where $\Lambda \equiv \mathcal{D} \circ \mathcal{E}$ with $\mathcal{E}: A \rightarrow T'_A Q$ and $\mathcal{D}: QB \rightarrow T'_B A' B'$ as defined above (for $n = 1$). Furthermore, set*

$$F := F(\Phi_{T'_A T'_B} \otimes \psi_{B' B R}, (\mathcal{D} \circ \mathcal{E} \otimes \text{id}_R)(\psi_{ABR})).$$

Then we have the following bounds for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:

$$\log F \leq \frac{1 - \alpha}{2\alpha} (\log |Q| - \log |T'_A| - S_\beta(R)_\rho + S_\alpha(AR)_\rho), \quad (3.30)$$

$$\log F \leq \frac{1 - \alpha}{2\alpha} (2 \log |Q| - S_\beta(R)_\rho + \tilde{S}_\alpha(R|A)_\rho). \quad (3.31)$$

Lemma 3.5 in turn implies the strong converse property for the quantum communication cost and entanglement gain of coherent state merging:

Theorem 3.6 (Strong converse for coherent state merging). *Let $\rho \equiv \rho_{AB}$ be a bipartite state and $\{(\rho^{\otimes n}, \Lambda_n)\}_{n \in \mathbb{N}}$ be a sequence of coherent state merging protocols as described above, with figure of merit F_n as defined in (3.28). Then for all $n \in \mathbb{N}$ we obtain the following bounds on the fidelity F_n for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:*

$$F_n \leq \exp \left\{ -n\kappa(\alpha) [S_\beta(AB)_\rho - S_\alpha(B)_\rho - q_{\text{csm}} + e_{\text{csm}}] \right\}, \quad (3.32)$$

$$F_n \leq \exp \left\{ -n\kappa(\alpha) [S_\beta(R)_\rho - \tilde{S}_\alpha(R|A)_\rho - 2q_{\text{csm}}] \right\}. \quad (3.33)$$

where $\kappa(\alpha) = (1 - \alpha)/(2\alpha)$, and $q_{\text{csm}} \equiv q_{\text{csm}}(\rho^{\otimes n}, \Lambda_n)$ and $e_{\text{csm}} \equiv e_{\text{csm}}(\rho^{\otimes n}, \Lambda_n)$ denote the quantum communication cost and entanglement gain respectively, as defined in (3.29).

3.6 Quantum state splitting

Quantum state splitting is the task in which a tripartite pure state ψ_{ACR} , which is initially shared between Alice (who has AC) and the reference (R), is split between Alice, Bob and the reference, with the system A being transferred to Bob. To this end, Alice and Bob can make use of prior shared entanglement and are allowed to do local operations on systems which they possess or receive. This protocol (also known as ‘Fully Quantum Reverse Shannon’ (FQRS) protocol [1, 15]) is dual to the coherent state merging protocol from Section 3.5 under time reversal [15]. Hence, the quantum state splitting protocol can also be obtained as a special case from the state redistribution protocol, if the systems B and T'_A are taken to be trivial. That is, Bob does not possess a share of the input state of the protocol, and the target state does not consist of an MES shared between Alice and Bob (i.e. the protocol always *consumes* entanglement).

Let Alice and Bob share n identical copies of the state ρ_{AC} with purification ψ_{ACR} , i.e. the state $\rho_{AC}^{\otimes n}$ with purification $\psi_{ACR}^{\otimes n}$. A general quantum state splitting protocol $(\rho^{\otimes n}, \Lambda_n)$ is

given by a joint quantum operation $\Lambda_n = \mathcal{D}_n \circ \mathcal{E}_n$ where $\mathcal{E}_n: A^n C^n T_A^n \rightarrow C^n Q^n$ is Alice's encoding map, the system Q^n is sent to Bob, and $\mathcal{D}_n: Q^n T_B^n \rightarrow A^n$ is Bob's decoding map. Here, $A^n \cong A^n$ and $C^n \cong C^n$. Denote the final state of the protocol by $\sigma_n = (\Lambda_n \otimes \text{id}_{R^n})(\Omega^n)$ where $\Omega^n \equiv \psi_{ACR}^{\otimes n} \otimes \Phi_{T_A^n T_B^n}^{k_n}$ is the initial state shared between Alice and Bob. Then the figure of merit is chosen to be the fidelity

$$F_n := F(\sigma_n, \psi^{\otimes n}), \quad (3.34)$$

where $\psi \equiv \psi_{A'C'R}$.

The *quantum communication cost* $q_{\text{qss}}(\rho^{\otimes n}, \Lambda_n)$ and the *entanglement cost* $e_{\text{qss}}(\rho^{\otimes n}, \Lambda_n)$ are defined in analogy to Section 3.1:

$$q_{\text{qss}}(\rho^{\otimes n}, \Lambda_n) := \frac{1}{n} \log |Q^n|, \quad e_{\text{qss}}(\rho^{\otimes n}, \Lambda_n) := \frac{1}{n} \log |T_A^n|. \quad (3.35)$$

A pair (e, q) , with $e, q \geq 0$, is said to be an achievable rate pair for quantum state splitting of a state ρ_{AC} , if there exists a sequence of protocols $\{(\rho^{\otimes n}, \Lambda_n)\}_{n \in \mathbb{N}}$ such that $\liminf_{n \rightarrow \infty} F_n = 1$ and

$$\limsup_{n \rightarrow \infty} e_{\text{qss}}(\rho^{\otimes n}, \Lambda_n) = e, \quad \limsup_{n \rightarrow \infty} q_{\text{qss}}(\rho^{\otimes n}, \Lambda_n) = q.$$

The optimal rates of entanglement cost and quantum communication cost for quantum state splitting were investigated in [1, 15, 5, 7]: A rate pair (e, q) is achievable if and only if e and q satisfy

$$q \geq \frac{1}{2} I(A; R)_\rho, \quad q + e \geq S(A)_\rho.$$

One-shot bounds characterizing the quantum communication cost and entanglement cost for quantum state splitting were derived by Berta *et al.* [7] as a building block in a proof of the Quantum Reverse Shannon theorem based on smooth entropies.

As mentioned at the beginning of this section, quantum state splitting is a special case of state redistribution with the choices $|B| = |T_A'| = 1$. In this case, $I(A; R|B)_\rho = I(A; R)$, and Lemma 3.1 reduces to

Lemma 3.7. *Let $\rho \equiv \rho_{AB}$ be a bipartite state with purification $|\psi_{ABR}\rangle$, and let (ρ, Λ) be a quantum state splitting protocol where $\Lambda \equiv \mathcal{D} \circ \mathcal{E}$ with $\mathcal{E}: AA'T_A \rightarrow AQ$ and $\mathcal{D}: QT_B \rightarrow B$ as defined above (for $n = 1$). Furthermore, set*

$$F := F(\psi_{ABR}, (\mathcal{D} \circ \mathcal{E} \otimes \text{id}_R)(\psi_{AA'R} \otimes \Phi_{T_A T_B})).$$

Then we have the following bounds for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:

$$\log F \leq \frac{1 - \alpha}{2\alpha} (\log |Q| + \log |T_A| - S_\beta(A)_\rho), \quad (3.36)$$

$$\log F \leq \frac{1 - \alpha}{2\alpha} (2 \log |Q| - S_\beta(R)_\rho + \tilde{S}_\alpha(R|A)_\rho), \quad (3.37)$$

$$\log F \leq \frac{1 - \alpha}{2\alpha} (2 \log |Q| - \tilde{I}_\alpha(R; A)_\rho). \quad (3.38)$$

Lemma 3.7 now implies a strong converse theorem for state splitting:

Theorem 3.8 (Strong converse for quantum state splitting). *Let $\rho \equiv \rho_{AB}$ be a bipartite state and $\{(\rho^{\otimes n}, \Lambda_n)\}_{n \in \mathbb{N}}$ be a sequence of quantum state splitting protocols as described above, with figure of merit F_n as defined in (3.34). Then for all $n \in \mathbb{N}$ we obtain the following bounds on the fidelity F_n for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:*

$$F_n \leq \exp \left\{ -n\kappa(\alpha) \left[S_\beta(A)_\rho - (q_{\text{qss}} + e_{\text{qss}}) \right] \right\}, \quad (3.39)$$

$$F_n \leq \exp \left\{ -n\kappa(\alpha) \left[S_\beta(R)_\rho - \tilde{S}_\alpha(R|A)_\rho - 2q_{\text{qss}} \right] \right\}, \quad (3.40)$$

$$F_n \leq \exp \left\{ -n\kappa(\alpha) \left[\tilde{I}_\alpha(R; A)_\rho - 2q_{\text{qss}} \right] \right\}, \quad (3.41)$$

where $\kappa(\alpha) = (1 - \alpha)/(2\alpha)$, and $q_{\text{qss}} \equiv q_{\text{qss}}(\rho^{\otimes n}, \Lambda_n)$ and $e_{\text{qss}} \equiv e_{\text{qss}}(\rho^{\otimes n}, \Lambda_n)$ denote the quantum communication cost and entanglement cost defined in (3.35), respectively.

4 Measurement compression with quantum side information

4.1 The protocol

Consider a bipartite state ρ_{AB} between two parties (say, Alice and Bob), and a positive operator-valued measure (POVM) $\Lambda = \{\Lambda_x\}_{x \in \mathcal{X}}$ (i.e. $0 \leq \Lambda_x \leq I_A$ for all $x \in \mathcal{X}$ and $\sum_{x \in \mathcal{X}} \Lambda_x = I_A$) on the A system, where X denotes a classical register. Suppose that Alice wants to communicate the measurement outcome X of Λ to Bob via classical communication. A simple solution is of course for Alice to apply the POVM Λ and send the outcome to Bob, requiring $\log |\mathcal{X}|$ bits of communication. In measurement compression with quantum side information, Alice and Bob want to reduce this communication cost by *simulating* the POVM Λ using shared randomness, Bob's quantum side information B , and sending $\log |L|$ bits of classical communication with $|L| \leq |\mathcal{X}|$. This information-theoretic task was introduced in [47] as an extension of Winter's original formulation of measurement compression [51]. In the following, we explain this protocol in more detail.

Given $\rho_{AB} \in \mathcal{D}(\mathcal{H}_{AB})$ and a POVM $\Lambda = \{\Lambda_x\}_{x \in \mathcal{X}}$ on the A system with outcome X , a general protocol for measurement compression with quantum side information consists of the following steps (see also Figure 5): Alice applies a quantum operation $\mathcal{E}: AM_A \rightarrow \tilde{X}L$ to her shares of a purification ψ_{RAB} of the initial state ρ_{AB} (with R being an inaccessible reference system) and the shared randomness $\chi_{M_A M_B}$. This produces a classical register \tilde{X} that holds her copy of the simulated outcome of the measurement, and a classical register L . She sends the latter to Bob, who then applies a quantum operation $\mathcal{D}: LBM_B \rightarrow \hat{X}B'$ to L and his shares of ψ_{RAB} and $\chi_{M_A M_B}$, producing a quantum output B' and the classical output \hat{X} , which represents the simulated outcome of the measurement. We denote the overall state of the protocol after applying \mathcal{E} and \mathcal{D} by ω and σ , respectively (cf. Figure 5). In comparison, we consider the ideal state $\varphi_{RXX'B}$ that would result from Alice applying the POVM Λ to her system A yielding the outcome X , and sending a copy X' uncompressed to Bob. That is, for $\zeta_A \in \mathcal{D}(\mathcal{H}_A)$ we define the measurement channel $\Lambda^{A \rightarrow XX'}(\zeta_A) := \sum_{x \in \mathcal{X}} \text{Tr}(\Lambda_x \zeta_A) |x\rangle\langle x|_X \otimes |x\rangle\langle x|_{X'}$ associated to the POVM Λ , and set $\varphi_{RXX'B} := (\text{id}_{RB} \otimes \Lambda^{A \rightarrow XX'}) (\psi_{RAB})$. The aim of the measurement compression protocol is to assure that σ is close in fidelity to the ideal state φ :

$$F := F(\varphi_{RXX'B}, \sigma_{R\tilde{X}\hat{X}B'}) = F(\varphi_{RXX'B}, (\text{id}_R \otimes \mathcal{D} \circ \mathcal{E})(\psi_{RAB} \otimes \chi_{M_A M_B})). \quad (4.1)$$

Given n identical copies of the input state ρ_{AB} with purification ψ_{RAB} and the shared randomness $\chi_{M_A M_B}^n$, we consider a measurement compression protocol with maps $\mathcal{E}_n: A^n M_A^n \rightarrow \tilde{X}^n L^n$ and $\mathcal{D}_n: L^n B^n M_B^n \rightarrow \hat{X}^n B'^n$, where L^n is the classical communication between Alice and

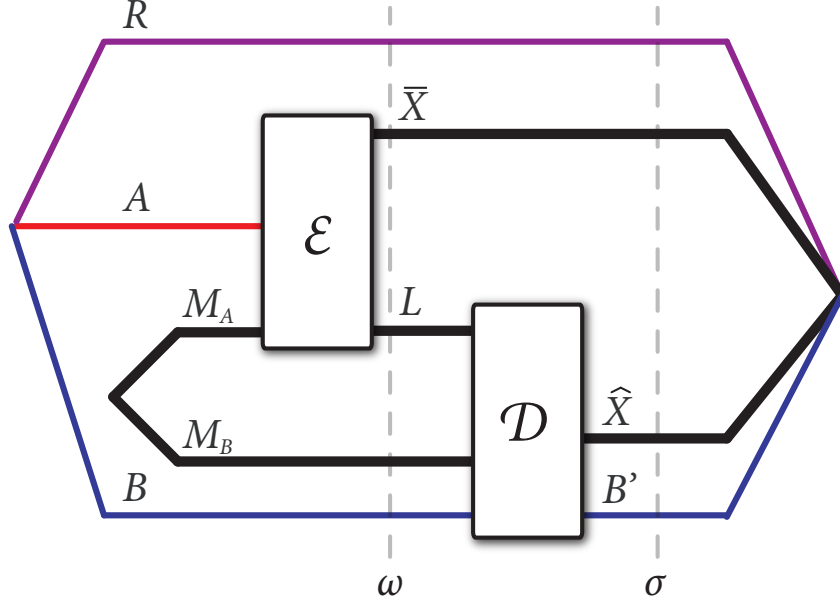


Figure 5: Measurement compression with quantum side information (see Section 4.1 for a detailed explanation).

Bob. The figure of merit is then given by

$$F_n := F \left(\varphi_n, (\text{id}_{R^n} \otimes \mathcal{D}_n \circ \mathcal{E}_n) \left(\psi_{RAB}^{\otimes n} \otimes \chi_{M_A^n M_B^n} \right) \right) \quad (4.2)$$

where the ideal state φ_n is obtained by Alice applying the POVM $\Lambda^{\otimes n}$ yielding the outcome X^n and sending a copy X^n to Bob. We define the *classical communication cost*

$$c(\rho^{\otimes n}, \Lambda^{\otimes n}) := \frac{1}{n} \log |L^n|, \quad (4.3)$$

and the *randomness cost*

$$r(\rho^{\otimes n}, \Lambda^{\otimes n}) := \frac{1}{n} \log |M_A^n|. \quad (4.4)$$

A rate pair (c, r) with $c, r \geq 0$ is called achievable if there exists a sequence $\{(\mathcal{E}_n, \mathcal{D}_n)\}_{n \in \mathbb{N}}$ of measurement compression protocols such that

$$\liminf_{n \rightarrow \infty} F_n = 1, \quad \limsup_{n \rightarrow \infty} c(\rho^{\otimes n}, \Lambda^{\otimes n}) = c, \quad \limsup_{n \rightarrow \infty} r(\rho^{\otimes n}, \Lambda^{\otimes n}) = r.$$

In [47] it was proved that (c, r) is achievable if and only if

$$c \geq I(X; R|B)_\varphi, \quad (4.5a)$$

$$c + r \geq S(X|B)_\varphi, \quad (4.5b)$$

where $\varphi_{RXX'B}$ is the ideal state of the protocol defined above.

4.2 Strong converse theorem

In this section, we strengthen the weak converse result obtained from (4.5a) for the classical communication cost in measurement compression with quantum side information to a strong converse theorem. As in Section 3.2, we first derive the following ‘one-shot’ lemma:

Lemma 4.1. *Let a bipartite state ρ_{AB} with purification ψ_{RAB} and a POVM Λ on A be given. Furthermore, let $\{(\mathcal{E}, \mathcal{D})\}$ be a measurement compression protocol as defined in Section 4.1 with figure of merit*

$$F := F(\varphi_{RXX'B'}, (\text{id}_R \otimes \mathcal{D} \circ \mathcal{E})(\psi_{RAB} \otimes \chi_{M_A M_B})).$$

Then we have the following bound on F for $\alpha \in (1/2, 1)$ and $\beta \equiv \beta(\alpha) = \alpha/(2\alpha - 1)$:

$$\log F \leq \frac{1-\alpha}{2\alpha} \left(\log |L| - \tilde{S}_\beta(R|B)_\varphi + \tilde{S}_\alpha(R|XB)_\varphi \right). \quad (4.6)$$

Proof. Define the states

$$\begin{aligned} \omega_{R\tilde{X}LBM_B} &:= (\text{id}_R \otimes \mathcal{E})(\psi_{RAB} \otimes \chi_{M_A M_B}), \\ \sigma_{R\tilde{X}\tilde{X}B'} &:= (\text{id}_R \otimes \mathcal{D} \circ \mathcal{E})(\psi_{RAB} \otimes \chi_{M_A M_B}). \end{aligned}$$

To prove (4.6), consider the following bound for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:

$$\begin{aligned} \frac{2\alpha}{1-\alpha} \log F &\leq \frac{2\alpha}{1-\alpha} \log F(\sigma_{R\tilde{X}B'}, \varphi_{RXB}), \\ &\leq \tilde{S}_\alpha(R|XB)_\varphi - \tilde{S}_\beta(R|\tilde{X}B')_\sigma \end{aligned} \quad (4.7)$$

where the first inequality follows from the monotonicity of the fidelity under partial trace, and the second inequality follows from eq. (2.16) of Proposition 2.8. We continue to bound the second term on the right-hand side of (4.7):

$$\begin{aligned} -\tilde{S}_\beta(R|\tilde{X}B')_\sigma &\leq -\tilde{S}_\beta(R|LBM_B)_\omega \\ &\leq \log |L| - \tilde{S}_\beta(R|BM_B)_\omega \\ &= \log |L| - \tilde{S}_\beta(R|BM_B)_{\psi \otimes \chi} \\ &= \log |L| - \tilde{S}_\beta(R|B)_\varphi. \end{aligned} \quad (4.8)$$

The first inequality follows from data processing with respect to the quantum operation $\mathcal{D}: LM_B \rightarrow \tilde{X}B'$ (Proposition 2.5(vii)), the second inequality follows from eq. (2.24) of Proposition 2.9, and the first equality follows from the fact that $\omega_{RBM_B} = \psi_{RB} \otimes \chi_{M_B}$. The last equality uses eq. (2.11) of Lemma 2.7, and the fact that $\psi_{RB} = \varphi_{RB}$. Substituting (4.8) in (4.7) then yields the claim. \square

This immediately implies the following strong converse theorem:

Theorem 4.2 (Strong converse theorem for measurement compression with QSI). *Let ρ_{AB} be a bipartite state, Λ a given POVM on A , and $\{(\mathcal{E}_n, \mathcal{D}_n)\}_{n \in \mathbb{N}}$ be a sequence of measurement compression protocols as described in Section 4.1, with figure of merit F_n as defined in (4.2). Then for all $n \in \mathbb{N}$ we have the following bound on F_n for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:*

$$F_n \leq \exp \left\{ -n\kappa(\alpha) \left[\tilde{S}_\beta(R|B)_\varphi - \tilde{S}_\alpha(R|XB)_\varphi - c \right] \right\}, \quad (4.9)$$

where $\kappa(\alpha) = (1 - \alpha)/(2\alpha)$ and $c \equiv c(\rho^{\otimes n}, \Lambda^{\otimes n})$ is the classical communication cost defined in (4.3).

Remark 4.3. The achievable rate region in the (c, r) -plane is determined by the two boundaries $c \geq I(X; R|B)_\varphi$ and $c + r \geq S(X|B)_\varphi$, as stated in (4.5) (compare this to the similar situation in the state redistribution protocol discussed in Section 3). Theorem 4.2 only proves the strong converse property for the c -boundary of the achievable rate region, and it remains open to prove

the strong converse property also for the $(c+r)$ -boundary as stated in (4.5b). While the proof of (4.6) in Lemma 4.1 closely follows that of (3.8) in Lemma 3.1, our investigations indicate that the proof method of (3.7) does not immediately carry over to show the desired bound for $c+r$ in measurement compression, that is,

$$\log F \stackrel{?}{\leq} f(\alpha)(\log |L| + \log |M_A| - S_{\beta(\alpha)}(XB)_\varphi + S_\alpha(B)_\varphi), \quad (4.10)$$

for some functions $f(\alpha)$ and $\beta(\alpha)$ satisfying $f(\alpha) > 0$ for all α in some open interval whose boundary contains 1, and $\lim_{\alpha \rightarrow 1} \beta(\alpha) = 1$.

5 Randomness extraction

5.1 The protocol

Suppose that Alice and Bob share the c-q state $\rho_{XB} = \sum_{x \in \mathcal{X}} p_x |x\rangle \langle x|_X \otimes \rho_B^x$ with $\rho_B^x \in \mathcal{D}(\mathcal{H}_B)$ for all $x \in \mathcal{X}$, where the classical register X is with Alice and the quantum system B is with Bob. The goal of a randomness extraction protocol is to extract from X a random string Z that is uncorrelated with B .

In the asymptotic, memoryless setting Alice and Bob share n copies of the c-q state ρ_{XB} . A general randomness extraction protocol $(\rho_{XB}^{\otimes n}, e_n)$ consists of a (surjective) encoding function $e_n: \mathcal{X}^n \rightarrow \mathcal{Z}^n$ where $\mathcal{X}^n = \mathcal{X}^{\times n}$. The classical register Z^n is then defined to be the one associated with the set \mathcal{Z}^n . The encoding function e_n gives rise to an encoding (quantum) operation, and without loss of generality this encoding map can be taken to be an isometry $U_{e_n}: |x^n\rangle \mapsto |x^n\rangle \otimes |e_n(x^n)\rangle$ where $|e_n(x^n)\rangle \in \{|z^n\rangle\}_{z^n \in \mathcal{Z}^n}$ for all $x^n \in \mathcal{X}^n$, resulting in the encoded state

$$\omega_{X^n Z^n B^n} := U_{e_n} \rho_{XB}^{\otimes n} U_{e_n}^\dagger = \sum_{x^n \in \mathcal{X}^n} p_{x^n} |x^n\rangle \langle x^n|_{X^n} \otimes |e_n(x^n)\rangle \langle e_n(x^n)|_{Z^n} \otimes \rho_{B^n}^{x^n}, \quad (5.1)$$

with $\rho_{B^n}^{x^n} := \rho_B^{x_1} \otimes \cdots \otimes \rho_B^{x_n}$ for $x^n = x_1 \dots x_n \in \mathcal{X}^n$. Upon discarding X^n , the final state of the protocol is then given by

$$\omega_n \equiv \omega_{Z^n B^n} = \sum_{z^n \in \mathcal{Z}^n} |z^n\rangle \langle z^n|_{Z^n} \otimes \sum_{x^n \in e_n^{-1}(z^n)} p_{x^n} \rho_{B^n}^{x^n}. \quad (5.2)$$

The randomness extraction protocol $(\rho_{XB}^{\otimes n}, e_n)$ succeeds if the final state ω_n is close to a state that is completely mixed on Z^n and independent of B^n . As the figure of merit we choose the fidelity

$$F_n := \max_{\sigma_{B^n}} F(\omega_n, \pi_{Z^n}^{\otimes n} \otimes \sigma_{B^n}). \quad (5.3)$$

The rate of extractable randomness $l(\rho_{XB}^{\otimes n}, e_n)$ is defined as

$$l(\rho_{XB}^{\otimes n}, e_n) := \frac{1}{n} \log |Z^n|. \quad (5.4)$$

A real number $l \geq 0$ is said to be an achievable rate for randomness extraction if there is a sequence of protocols $\{(\rho_{XB}^{\otimes n}, e_n)\}_{n \in \mathbb{N}}$ such that

$$\liminf_{n \rightarrow \infty} F_n = 1, \quad \liminf_{n \rightarrow \infty} l(\rho_{XB}^{\otimes n}, e_n) = l.$$

Randomness extraction was first studied by Bennett *et al.* [4] (under the name of ‘privacy amplification’) and further developed by Renner [38] and Renner and König [39] (see also [45]). They showed that $l \geq 0$ is an achievable rate for randomness extraction if and only if

$$l \leq S(X|B)_\rho. \quad (5.5)$$

Tomamichel [42] proved a strong converse theorem for randomness extraction based on one-shot bounds in terms of smooth entropies.

5.2 Strong converse theorem

We first state the following general bound on the fidelity:

Lemma 5.1 ([41]). *Let $\rho_{AB} \in \mathcal{D}(\mathcal{H}_{AB})$, $\sigma_A \in \mathcal{D}(\mathcal{H}_A)$, and $\chi_B \in \mathcal{D}(\mathcal{H}_B)$ be arbitrary quantum states, then*

$$F(\rho_{AB}, \sigma_A \otimes \rho_B) \geq F^2(\rho_{AB}, \sigma_A \otimes \chi_B).$$

With this result in hand, we can prove the following bound on the fidelity for arbitrary randomness extraction protocols:

Lemma 5.2. *Let $\rho_{XB} = \sum_{x \in \mathcal{X}} p_x |x\rangle\langle x|_X \otimes \rho_B^x$ be a c - q state with $\rho_B^x \in \mathcal{D}(\mathcal{H}_B)$ for all $x \in \mathcal{X}$, and denote by ω_{XZB} the encoded state of a randomness extraction protocol (ρ_{XB}, e) as defined in Section 5.1 (for $n = 1$). Furthermore, set*

$$F := \max_{\sigma_B} F(\omega_{ZB}, \pi_Z \otimes \sigma_B).$$

Then we have the following bound for all $\alpha \in (1/2, 1)$ and $\beta \equiv \beta(\alpha) = \alpha / (2\alpha - 1)$:

$$\log F \leq \frac{1 - \alpha}{4\alpha} (S_\alpha(XB)_\rho - S_\beta(B)_\rho - \log |Z|). \quad (5.6)$$

For the same range of α , we also obtain the following alternative bound on F :

$$\log F \leq \frac{1 - \alpha}{4\alpha} (\tilde{S}_\alpha(X|B)_\rho - \log |Z|). \quad (5.7)$$

Proof. We first prove (5.6). Set $F' := F(\omega_{ZB}, \pi_Z \otimes \rho_B)$, then by eq. (2.15) of Proposition 2.8 we have the following bound for $\alpha \in (1/2, 1)$ and $\beta = \alpha / (2\alpha - 1)$:

$$\begin{aligned} S_\alpha(ZB)_\omega &\geq S_\beta(ZB)_{\pi_Z \otimes \rho} + \frac{2\alpha}{1 - \alpha} \log F' \\ &= \log |Z| + S_\beta(B)_\rho + \frac{2\alpha}{1 - \alpha} \log F'. \end{aligned} \quad (5.8)$$

In the second line we used additivity of the Rényi entropy, as well as the fact that $S_\gamma(\pi_Z) = \log |Z|$ for all $\gamma \geq 0$ (cf. Proposition 2.5(iii) and (ii)).

Furthermore, we have the bound

$$S_\alpha(XB)_\rho = S_\alpha(XZB)_\omega \geq S_\alpha(ZB)_\omega, \quad (5.9)$$

where the equality follows from the invariance of the Rényi entropy under the encoding isometry U_e (Proposition 2.5(iv)), and the inequality follows from eq. (2.23) of Proposition 2.9.

Putting (5.8) and (5.9) together, we obtain

$$\log F' \leq \frac{1-\alpha}{2\alpha} (S_\alpha(XB)_\rho - S_\beta(B)_\rho - \log |Z|). \quad (5.10)$$

Now observe that $\omega_B = \rho_B$, and by Lemma 5.1 we have

$$F' = F(\omega_{ZB}, \pi_Z \otimes \rho_B) = F(\omega_{ZB}, \pi_Z \otimes \omega_B) \geq F^2(\omega_{ZB}, \pi_Z \otimes \sigma_B) \quad (5.11)$$

for all $\sigma_B \in \mathcal{D}(\mathcal{H}_B)$. Substituting this into (5.10) and using the monotonicity of the logarithm then yields the claim.

To prove (5.7), observe that (5.11) together with eq. (2.16) of Proposition 2.8 yield the following for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:

$$\begin{aligned} \frac{4\alpha}{1-\alpha} \log F &\leq \frac{2\alpha}{1-\alpha} \log F' \\ &\leq \tilde{S}_\alpha(Z|B)_\omega - \tilde{S}_\beta(Z|B)_{\pi \otimes \rho} \\ &= \tilde{S}_\alpha(Z|B)_\omega - \log |Z| \\ &\leq \tilde{S}_\alpha(XZ|B)_\omega - \log |Z| \\ &= \tilde{S}_\alpha(X|B)_\rho - \log |Z|, \end{aligned}$$

where the first equality follows from eq. (2.11) of Lemma 2.7, the third inequality uses eq. (2.23) of Proposition 2.9, and the last equality uses the invariance of the Rényi conditional entropy under the isometry U_e . \square

This implies the following strong converse theorem for randomness extraction:

Theorem 5.3. *Let ρ_{XB} be a c - q state, and let $\{(\rho_{XB}^{\otimes n}, e_n)\}_{n \in \mathbb{N}}$ be a sequence of randomness extraction protocols as defined in Section 5.1 with figure of merit F_n as given by (5.3). Then for all $n \in \mathbb{N}$ we have the following bounds on F_n for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:*

$$F_n \leq \exp \left\{ -\kappa(\alpha)n \left[l - S_\alpha(XB)_\rho + S_\beta(B)_\rho \right] \right\}, \quad (5.12)$$

$$F_n \leq \exp \left\{ -\kappa(\alpha)n \left[l - \tilde{S}_\alpha(X|B)_\rho \right] \right\}, \quad (5.13)$$

where $\kappa(\alpha) = (1 - \alpha)/(4\alpha)$, and $l \equiv l(\rho_{XB}^{\otimes n}, e_n)$ is the rate of extractable randomness defined in (5.4).

Remark 5.4. An immediate question arising from Theorem 5.3 is whether one of the two bounds in (5.12) and (5.13) is tighter than the other, that is, whether one of

$$S_\alpha(XB)_\rho - S_\beta(B)_\rho \stackrel{?}{\leq} \tilde{S}_\alpha(X|B)_\rho, \quad S_\alpha(XB)_\rho - S_\beta(B)_\rho \stackrel{?}{\geq} \tilde{S}_\alpha(X|B)_\rho \quad (5.14)$$

holds for all $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$. Numerical investigations with classical registers B show that neither inequality in (5.14) is always valid. Hence, the two exponents in (5.12) and (5.13) are in general incomparable. In the light of identifying the strong converse exponent of randomness extraction against quantum side information (cf. Section 7), this fact indicates that further analysis is needed here, as a strong converse exponent usually characterizes the tightest possible strong converse bound. Interchanging α and β in the above, the same reasoning applies to Theorem 6.2 of the following section.

6 Data compression with quantum side information

6.1 The protocol

In data compression with quantum side information, Alice has a classical register X , and Bob holds the quantum system B (the ‘side information’) which is correlated with X . The goal of the protocol is for Alice to encode her classical register X in a (smaller) classical register C such that Bob can recover X from C and his quantum system B . A data compression protocol $(\rho_{XB}, e, \mathcal{D})$ can be described by the following steps:

The initial state is a classical-quantum state $\rho_{XB} = \sum_{x \in \mathcal{X}} p_x |x\rangle\langle x|_X \otimes \rho_B^x$ where the classical register X is with Alice and the quantum system B is with Bob. To encode her message in the system C , Alice uses an arbitrary encoding function $e : \mathcal{X} \rightarrow \mathcal{C}$ (which we assume to be surjective, i.e. $\mathcal{C} = e(\mathcal{X})$). The classical register C is then defined as the one associated with the Hilbert space \mathcal{H}_C with orthonormal basis $\{|c\rangle\}_{c \in \mathcal{C}}$. The encoding function e gives rise to an encoding (quantum) map, which without loss of generality can be taken to be an isometric encoding map $U_e : |x\rangle \mapsto |x\rangle \otimes |e(x)\rangle$ where $|e(x)\rangle \in \{|c\rangle\}_{c \in \mathcal{C}}$ for all $x \in \mathcal{X}$. This results in the state

$$\omega_{XCB} := U_e \rho_{XB} U_e^\dagger = \sum_{x \in \mathcal{X}} p_x |x\rangle\langle x|_X \otimes |e(x)\rangle\langle e(x)|_C \otimes \rho_B^x. \quad (6.1)$$

Upon receiving the classical message C , Bob applies a measurement given by a POVM $\Lambda_c = \{\Lambda_{x',c}\}_{x' \in \mathcal{X}}$ to his state ρ_B^x , where the measurement is conditioned on the value of c in the encoded register C . We label the corresponding random variable by X' . The final state of the classical registers X and X' is given by

$$\sigma_{XX'} := \sum_{x, x' \in \mathcal{X}} p_x q_{x'|x} |xx'\rangle\langle xx'|_{XX'} \quad \text{where} \quad q_{x'|x} := \text{Tr}(\Lambda_{x',e(x)} \rho_B^x). \quad (6.2)$$

The POVM constitutes a quantum operation $\mathcal{D} : CB \rightarrow X'$, that is,

$$\sigma_{XX'} = (\text{id}_X \otimes \mathcal{D})(\omega_{XCB}).$$

Here, the CPTP map \mathcal{D} is defined as the one that, conditioned on the value c , applies the map

$$\mathcal{D}^c(v_B) = \sum_{x' \in \mathcal{X}} \text{Tr}(\Lambda_{x',c} v_B) |x'\rangle\langle x'|_{X'} \quad \text{for } v_B \in \mathcal{D}(\mathcal{H}_B), \quad (6.3)$$

implementing the POVM Λ_c to the B system, and then traces out C . Note that \mathcal{D}^c is a special case of an entanglement-breaking channel [25] with Kraus operators given by

$$K_{x'j}^c := |x'\rangle\langle j_B| \sqrt{\Lambda_{x',c}} \quad (6.4)$$

for an orthonormal basis $\{|j_B\rangle\}_{j=1}^{|B|}$ of \mathcal{H}_B .

The final state $\sigma_{XX'}$ of the protocol should be close to the classically correlated target state $\varphi_{XX'}$ defined by

$$\varphi_{XX'} := \sum_{x \in \mathcal{X}} p_x |xx\rangle\langle xx|_{XX'}. \quad (6.5)$$

As the figure of merit for a data compression protocol $(\rho_{XB}, e, \mathcal{D})$ we choose the success probability $p_{\text{succ}}(\rho_{XB}, e, \mathcal{D})$ of successfully decoding X :

$$p_{\text{succ}}(\rho_{XB}, e, \mathcal{D}) := \sum_{x \in \mathcal{X}} p_x q_{x|x}. \quad (6.6)$$

If Alice and Bob share n copies of the c-q state ρ_{XB} , then the figure of merit for a data compression protocol $(\rho_{XB}^{\otimes n}, e_n, \mathcal{D}_n)$ is given by

$$p_n := p_{\text{succ}}(\rho_{XB}^{\otimes n}, e_n, \mathcal{D}_n) = \sum_{x^n \in \mathcal{X}^n} p_{x^n} q_{x^n | x^n}, \quad (6.7)$$

where $e_n: \mathcal{X}^n \rightarrow \mathcal{C}^n$ is the encoding function, and $p_{x^n} q_{x^n | x^n}$ is the probability distribution of the classical state

$$\sigma_{X^n X^n} := (\text{id}_{X^n} \otimes \mathcal{D}_n) \left(U_{e_n} \rho_{XB}^{\otimes n} U_{e_n}^\dagger \right).$$

Here, the CPTP map \mathcal{D}_n implements the POVM on the system B^n . The rate $m(\rho_{XB}^{\otimes n}, e_n, \mathcal{D}_n)$ of the data compression protocol $(\rho_{XB}^{\otimes n}, e_n, \mathcal{D}_n)$ is defined as

$$m(\rho_{XB}^{\otimes n}, e_n, \mathcal{D}_n) := \frac{1}{n} \log |\mathcal{C}^n|. \quad (6.8)$$

A real number $m \geq 0$ is said to be an achievable rate for data compression with quantum side information, if there is a sequence $\{(\rho_{XB}^{\otimes n}, e_n, \mathcal{D}_n)\}_{n \in \mathbb{N}}$ of protocols satisfying

$$\limsup_{n \rightarrow \infty} m(\rho_{XB}^{\otimes n}, e_n, \mathcal{D}_n) = m, \quad \liminf_{n \rightarrow \infty} p_n = 1.$$

Data compression with quantum side information is the ‘dual’ task [36] to randomness extraction, with the optimal rate of the former also given by the conditional entropy $S(X|B)_\rho$. That is, a real number $m \geq 0$ is an achievable rate for data compression with quantum side information if and only if

$$m \geq S(X|B)_\rho.$$

This was proved by Winter [49] (see also Devetak and Winter [16]). Renes and Renner [37] derived one-shot bounds for data compression with quantum side information in terms of smooth entropies. Tomamichel [42] proved a strong converse theorem for data compression based on one-shot bounds in terms of smooth entropies.

6.2 Strong converse theorem

Lemma 6.1. *Let $\rho_{XB} = \sum_{x \in \mathcal{X}} p_x |x\rangle\langle x|_X \otimes \rho_B^x$ be a c-q state with $\rho_B^x \in \mathcal{D}(\mathcal{H}_B)$ for all $x \in \mathcal{X}$, and consider a data compression protocol $(\rho_{XB}, e, \mathcal{D})$ as defined in Section 6.1. We have the following bounds for the success probability $p_{\text{succ}} \equiv p_{\text{succ}}(\rho_{XB}, e, \mathcal{D})$ defined in (6.6), for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:*

$$\log p_{\text{succ}} \leq \frac{1 - \alpha}{2\alpha} (\log |\mathcal{C}| - S_\beta(XB)_\rho + S_\alpha(B)_\rho), \quad (6.9)$$

$$\log p_{\text{succ}} \leq \frac{1 - \alpha}{2\alpha} (\log |\mathcal{C}| - \tilde{S}_\beta(X|B)_\rho). \quad (6.10)$$

Proof. We first prove (6.9). Given an arbitrary POVM $\Lambda_c = \{\Lambda_{x',c}\}_{x' \in \mathcal{X}}$ for $c \in \mathcal{C}$ and the corresponding conditional decoding map \mathcal{D}^c defined in (6.3) with Kraus operators $K_{x',j}^c$ given by (6.4), it is straightforward to construct from this a Stinespring isometry $V_c: \mathcal{H}_B \rightarrow \mathcal{H}_{X'} \otimes \mathcal{H}_E$ for \mathcal{D}^c :

$$V_c := \sum_{x',j} K_{x',j}^c \otimes |x'j\rangle_E = \sum_{x',j} |x'j\rangle_{X'E} \langle j_B | \Lambda_{x',c}^{1/2} \otimes |x'j\rangle_E \quad (6.11)$$

where $\{|x'j_E\rangle\}_{x' \in \mathcal{X}, j=1, \dots, |B|}$ is an orthonormal basis for the environment \mathcal{H}_E with $\dim \mathcal{H}_E = |\mathcal{X}||B|$, satisfying $\langle xj|yk\rangle_E = \delta_{xy}\delta_{jk}$. The isometry defined in (6.11) satisfies $V_c^\dagger V_c = I_B$ and $\mathcal{D}_c(\rho_B) = \text{Tr}_E(V_c \rho_B V_c^\dagger)$.

For every $x \in \mathcal{X}$, let $|\psi_{BS}^x\rangle$ be a purification of ρ_B^x . Consider then the following purification of ρ_{XB} :

$$|\psi_{XBRs}\rangle := \sum_{x \in \mathcal{X}} \sqrt{p_x} |x_X\rangle |x_R\rangle |\psi_{BS}^x\rangle.$$

Then the pure states obtained after encoding with U_e and decoding with \mathcal{D} in the data compression protocol (cf. Section 6.1) are given by:

$$\begin{aligned} |\omega_{XCBS}\rangle &= U_e |\psi_{XBRs}\rangle \\ &= \sum_{x \in \mathcal{X}} \sqrt{p_x} |x_X\rangle |x_R\rangle |e(x)_C\rangle |\psi_{BS}^x\rangle, \end{aligned} \quad (6.12a)$$

$$\begin{aligned} |\sigma_{XX'RSE}\rangle &= V |\omega_{XCBS}\rangle \\ &= \sum_{x, x' \in \mathcal{X}} \sum_{j=1, \dots, |B|} \sqrt{p_x} |x_X\rangle |x_R\rangle |x'_{X'}\rangle \langle j_B | \Lambda_{x', e(x)}^{1/2} | \psi_{BS}^x \rangle |x'j_E\rangle. \end{aligned} \quad (6.12b)$$

Here, V denotes the Stinespring isometry of the overall decoding map \mathcal{D} that, conditional on the classical message $c \in \mathcal{C}$, applies the decoding operation \mathcal{D}^c . It can easily be checked that $|\omega_{XCBS}\rangle$ and $|\sigma_{XX'RSE}\rangle$, as given by (6.12a) and (6.12b) respectively, indeed purify ω_{XCB} and $\sigma_{XX'}$ as given by (6.1) and (6.2).

In the next step, we relate the final state $\sigma_{XX'}$ (resp. its purification $|\sigma_{XX'RSE}\rangle$) of the data compression protocol to the target state $\varphi_{XX'}$ given by (6.5), thus obtaining a bound on the success probability $p_{\text{succ}} \equiv p_{\text{succ}}(\rho_{XB}, e, \mathcal{D})$ defined in (6.6). To this end, we consider the following data compression protocol that allows perfect recovery of the register X , resulting in the target state $\varphi_{XX'}$: Take $\mathcal{C} \cong X$ and consider the trivial encoding $e(x) = x$ for all $x \in \mathcal{X}$, together with the trivial POVM $E_c = \{\delta_{x',c} I_B\}_{x'}$ which discards the quantum system B and yields the message $x' = c = x$ with certainty. Following the same procedure as above, we can work out the pure states obtained after encoding and decoding using these particular choices:

$$|\bar{\omega}_{XCBS}\rangle = \sum_{x \in \mathcal{X}} \sqrt{p_x} |x_X\rangle |x_R\rangle |x_C\rangle |\psi_{BS}^x\rangle, \quad (6.13a)$$

$$|\bar{\sigma}_{XX'RSE}\rangle = \sum_{x \in \mathcal{X}} \sum_{j=1, \dots, |B|} \sqrt{p_x} |x_X\rangle |x_R\rangle |x_{X'}\rangle \langle j_B | \psi_{BS}^x \rangle |xj_E\rangle. \quad (6.13b)$$

Note that $\bar{\sigma}_{XX'RSE}$ indeed purifies the target state $\varphi_{XX'}$ of the data compression protocol. Let us compute the inner product of the pure states σ and $\bar{\sigma}$:

$$\begin{aligned} \langle \bar{\sigma}_{XX'RSE} | \sigma_{XX'RSE} \rangle &= \sum_{x, y, y' \in \mathcal{X}} \sum_{j, k} \sqrt{p_x p_y} \langle xx | yy \rangle_{XR} \langle x | y' \rangle_{X'} \langle \psi_{BS}^x | j \rangle \langle k | \Lambda_{y', e(y)}^{1/2} | \psi_{BS}^x \rangle \langle xj | y'k \rangle_E \\ &= \sum_{x \in \mathcal{X}} \sum_j p_x \langle \psi_{BS}^x | j \rangle \langle j | \Lambda_{x, e(x)}^{1/2} | \psi_{BS}^x \rangle \\ &= \sum_{x \in \mathcal{X}} p_x \langle \psi_{BS}^x | \Lambda_{x, e(x)}^{1/2} | \psi_{BS}^x \rangle \\ &\geq \sum_{x \in \mathcal{X}} p_x \langle \psi_{BS}^x | \Lambda_{x, e(x)} | \psi_{BS}^x \rangle \\ &= \sum_{x \in \mathcal{X}} p_x \text{Tr} \left(\Lambda_{x, e(x)} \rho_B^x \right) \\ &= p_{\text{succ}}, \end{aligned} \quad (6.14)$$

where in the third equality we used the completeness of the basis $\{|j_B\rangle\}_{j=1}^{|B|}$, and in the inequality we used the fact that

$$\sqrt{\Lambda_{x',c}} \geq \Lambda_{x',c} \quad \text{for all } x' \in \mathcal{X} \text{ and } c \in \mathcal{C},$$

since $0 \leq \Lambda_{x',c} \leq I_B$ for all $x' \in \mathcal{X}$ and $c \in \mathcal{C}$. On the other hand, we have

$$\begin{aligned} \langle \bar{\sigma}_{XX'RSE} | \sigma_{XX'RSE} \rangle &= |\langle \bar{\sigma}_{XX'RSE} | \sigma_{XX'RSE} \rangle| \\ &= F(\bar{\sigma}_{XX'RSE}, \sigma_{XX'RSE}) \\ &\leq F(\bar{\sigma}_{XRS}, \sigma_{XRS}) \\ &= F(\bar{\sigma}_{XRS}, \omega_{XRS}), \end{aligned} \tag{6.15}$$

where we used the monotonicity of the fidelity under partial trace in the inequality, and the last line follows from the fact that the decoding \mathcal{D} does not affect the systems X , R , and S . Putting (6.14) and (6.15) together, we obtain the following bound from eq. (2.15) of Proposition 2.8 for $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$:

$$S_\alpha(XRS)_\omega \geq S_\beta(XRS)_{\bar{\sigma}} + \frac{2\alpha}{1-\alpha} \log p_{\text{succ}}. \tag{6.16}$$

By the duality property of the Rényi entropies (Proposition 2.5(v)) and the subadditivity property (Lemma 2.6), we have

$$S_\alpha(XRS)_\omega = S_\alpha(CB)_\omega \leq \log |C| + S_\alpha(B)_\omega = \log |C| + S_\alpha(B)_\rho, \tag{6.17}$$

using the fact that $\omega_B = \rho_B$. Furthermore, due to the choice of the trivial encoding and decoding operations defined above and resulting in the state $\bar{\sigma}_{XX'} = \varphi_{XX'}$, we have

$$S_\beta(XRS)_{\bar{\sigma}} = S_\beta(X'E)_{\bar{\sigma}} = S_\beta(CB)_{\bar{\omega}} = S_\beta(XB)_\rho, \tag{6.18}$$

where we used the invariance of the Rényi entropies under isometries (Proposition 2.5(iv)) in the second equality, and the fact that $C \cong X$ is just a copy of the initial register X in the last equality. Hence, combining (6.16), (6.17), and (6.18) we obtain

$$\log |C| + S_\alpha(B)_\rho \geq S_\beta(XB)_\rho + \frac{2\alpha}{1-\alpha} \log p_{\text{succ}},$$

which after re-arranging yields (6.9).

To prove (6.10), consider the following chain of inequalities:

$$\begin{aligned} \frac{2\alpha}{1-\alpha} \log p_{\text{succ}} &\leq \frac{2\alpha}{1-\alpha} \log F(\sigma_{XX'}, \varphi_{XX'}) \\ &\leq \tilde{S}_\alpha(X|X')_\varphi - \tilde{S}_\beta(X|X')_\sigma \\ &= -\tilde{S}_\beta(X|X')_\sigma \\ &\leq -\tilde{S}_\beta(X|CB)_\omega \\ &\leq \log |C| - \tilde{S}_\beta(XC|B)_\omega \\ &= \log |C| - \tilde{S}_\beta(X|B)_\rho. \end{aligned}$$

The first inequality follows from (6.14) and (6.15), the second inequality uses eq. (2.16) of Proposition 2.8, and the first equality uses the fact that $\tilde{S}_\alpha(X|X')_\varphi = 0$. The third inequality is data processing with respect to the decoding map \mathcal{D} , the fourth inequality follows from [30, Prop. 8], and the last equality uses the invariance of the Rényi conditional entropy with respect to the encoding isometry U_e . \square

We then have the following strong converse theorem for data compression with quantum side information:

Theorem 6.2. Let ρ_{XB} be a c - q state and let $\{(\rho_{XB}^{\otimes n}, e_n, \mathcal{D}_n)\}_{n \in \mathbb{N}}$ be a sequence of data compression protocols as defined in Section 6.1, with figure of merit $p_n \equiv p_{\text{succ}}(\rho_{XB}^{\otimes n}, e_n, \mathcal{D}_n)$ as defined in (6.7). Then for all $n \in \mathbb{N}$ we have the following bounds on p_n for $\alpha \in (1/2, 1)$ and $\beta = \alpha / (2\alpha - 1)$:

$$p_n \leq \exp \left\{ -n\kappa(\alpha) \left[S_\beta(XB)_\rho - S_\alpha(B)_\rho - m \right] \right\}, \quad (6.19)$$

$$p_n \leq \exp \left\{ -n\kappa(\alpha) \left[\tilde{S}_\beta(X|B)_\rho - m \right] \right\}, \quad (6.20)$$

where $\kappa(\alpha) = (1 - \alpha) / (2\alpha)$, and $m \equiv m(\rho_{XB}^{\otimes n}, e_n, \mathcal{D}_n)$ denotes the compression rate as defined in (6.8).

Remark 6.3. See Remark 5.4 for a discussion of whether the two strong converse bounds of Theorem 6.2 are comparable.

7 Discussion and open questions

For any information-theoretic task one can define *achievable* and *strong converse rates*. An achievable rate is a non-negative real number such that if one codes at a rate above it (for the case in which the optimal rate is a cost) then the error probability of the protocol vanishes asymptotically. The strong converse rate, in contrast, is a non-negative real number such that if one codes at a rate below it, then the error probability goes to one in the asymptotic limit. Consequently, the fidelity between the initial and final target states of the protocol decays to zero asymptotically. The optimal rate of the protocol is said to satisfy the *strong converse property* if the largest strong converse rate coincides with the smallest achievable rate. In this case the optimal rate provides a sharp rate threshold for the task. The exact exponent of the decaying fidelity (or success probability) for a given rate above zero and below the smallest achievable rate is called the *strong converse exponent*.⁷

In this paper, we used the Rényi entropy method to derive strong converse theorems which establish the strong converse property of the optimal rates for the following protocols: state redistribution (without and with feedback) (Theorem 3.2 and Theorem 3.4), coherent state merging (Theorem 3.6), quantum state splitting (Theorem 3.8), measurement compression with quantum side information (Theorem 4.2), randomness extraction against quantum side information (Theorem 5.3), and data compression with quantum side information (Theorem 6.2).

To this end, we established certain Rényi generalizations of the optimal rates of these protocols as bounds on the strong converse exponents. We derived inequalities involving these Rényi generalizations, the most important of which relate Rényi entropic quantities to the fidelity between two quantum states. These inequalities played a crucial role in the proofs of our strong converse theorems.

Determining whether any of these Rényi generalizations are indeed the exact strong converse exponents for the information-theoretic tasks in question is an interesting problem for future research.

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⁷If the optimal rate is a *gain* instead of a cost, one needs to swap ‘above’ and ‘below’ as well as ‘smallest’ and ‘largest’ in the above paragraph.

A Remaining proofs of Section 2.4

A.1 Lemma 2.7

Proof of (2.10) in Lemma 2.7. Consider that

$$\begin{aligned}\tilde{I}_\alpha(A; CB)_\rho &= \min_{\tau_{CB}} \frac{\alpha}{\alpha-1} \log \left\| \left(\rho_A^{(1-\alpha)/2\alpha} \otimes \tau_{CB}^{(1-\alpha)/2\alpha} \right) \rho_{ABC} \left(\rho_A^{(1-\alpha)/2\alpha} \otimes \tau_{CB}^{(1-\alpha)/2\alpha} \right) \right\|_\alpha \\ &= \min_{\tau_{CB}} \frac{\alpha}{\alpha-1} \log \left\| \tau_{CB}^{(1-\alpha)/2\alpha} \tilde{\rho}_{ABC} \tau_{CB}^{(1-\alpha)/2\alpha} \right\|_\alpha + \frac{\alpha}{\alpha-1} \log \text{Tr} \left\{ \rho_A^{(1-\alpha)/\alpha} \rho_{AB} \right\} \\ &= -\tilde{S}_\alpha(A|CB)_{\tilde{\rho}} + \frac{\alpha}{\alpha-1} \log \text{Tr} \left\{ \rho_A^{(1-\alpha)/\alpha} \rho_{AB} \right\},\end{aligned}$$

where we have defined the density operator

$$\tilde{\rho}_{ABC} \equiv \frac{1}{\text{Tr} \left\{ \rho_A^{(1-\alpha)/\alpha} \rho_{ABC} \right\}} \rho_A^{(1-\alpha)/2\alpha} \rho_{ABC} \rho_A^{(1-\alpha)/2\alpha}$$

and have observed that

$$\text{Tr} \left\{ \rho_A^{(1-\alpha)/\alpha} \rho_{ABC} \right\} = \text{Tr} \left\{ \rho_A^{(1-\alpha)/\alpha} \rho_{AB} \right\}.$$

Furthermore,

$$\tilde{\rho}_{AB} = \text{Tr}_C \left\{ \tilde{\rho}_{ABC} \right\} = \frac{1}{\text{Tr} \left\{ \rho_A^{(1-\alpha)/\alpha} \rho_{AB} \right\}} \rho_A^{(1-\alpha)/2\alpha} \rho_{AB} \rho_A^{(1-\alpha)/2\alpha}.$$

We now apply the bound from Lemma 2.7, eq. (2.9),

$$-\tilde{S}_\alpha(A|CB)_{\tilde{\rho}} \leq -\tilde{S}_\alpha(A|B)_{\tilde{\rho}} + 2 \log |C|,$$

to see that

$$\begin{aligned}-\tilde{S}_\alpha(A|CB)_{\tilde{\rho}} + \frac{\alpha}{\alpha-1} \log \text{Tr} \left\{ \rho_A^{(1-\alpha)/\alpha} \rho_{AB} \right\} \\ \leq -\tilde{S}_\alpha(A|B)_{\tilde{\rho}} + \frac{\alpha}{\alpha-1} \log \text{Tr} \left\{ \rho_A^{(1-\alpha)/\alpha} \rho_{AB} \right\} + 2 \log |C| \\ = \tilde{I}_\alpha(A; Bt)_\rho + 2 \log |C|,\end{aligned}$$

which yields (2.10). \square

Proof of (2.12) in Lemma 2.7. From the data processing inequality (Proposition 2.5(viii)) we know that

$$\tilde{I}_\alpha(A; BC)_{\rho \otimes \sigma} \geq \tilde{I}_\alpha(A; B)_\rho. \quad (\text{A.1})$$

On the other hand, consider that

$$\tilde{I}_\alpha(A; BC)_{\rho \otimes \sigma} = \min_{\tau_{BC}} \tilde{D}_\alpha(\rho_{AB} \otimes \sigma_C \| \rho_A \otimes \tau_{BC}) \quad (\text{A.2})$$

$$\leq \tilde{D}_\alpha(\rho_{AB} \otimes \sigma_C \| \rho_A \otimes \theta_B \otimes \sigma_C) \quad (\text{A.3})$$

$$= \tilde{D}_\alpha(\rho_{AB} \| \rho_A \otimes \theta_B). \quad (\text{A.4})$$

Since the inequality holds for all θ_B , we get that

$$\tilde{I}_\alpha(A; BC)_{\rho \otimes \sigma} \leq \tilde{I}_\alpha(A; B)_\rho, \quad (\text{A.5})$$

and we are done. \square

A.2 Proposition 2.8

Proof of (2.17) in Proposition 2.8. For an arbitrary density operator τ_B and $\varepsilon \in (0, 1)$ define the state $\tau(\varepsilon)_B := (1 - \varepsilon)\tau_B + \varepsilon\pi_B$. We then have the following chain of inequalities:

$$\begin{aligned}
& \tilde{D}_\beta(\rho_{AB} \parallel \rho_A \otimes \tau_B) - \tilde{D}_\alpha(\sigma_{AB} \parallel \sigma_A \otimes \tau(\varepsilon)_B) - \log(1 - \varepsilon) \\
&= \tilde{D}_\beta(\rho_{AB} \parallel \rho_A \otimes (1 - \varepsilon)\tau_B) - \tilde{D}_\alpha(\sigma_{AB} \parallel \sigma_A \otimes \tau(\varepsilon)_B) \\
&\geq \tilde{D}_\beta(\rho_{AB} \parallel \rho_A \otimes \tau(\varepsilon)_B) - \tilde{D}_\alpha(\sigma_{AB} \parallel \sigma_A \otimes \tau(\varepsilon)_B) \\
&= \frac{2\beta}{\beta - 1} \log \left\| \rho_{AB}^{1/2} \rho_A^{(1-\beta)/2\beta} \otimes \tau(\varepsilon)_B^{(1-\beta)/2\beta} \right\|_{2\beta} - \frac{2\alpha}{\alpha - 1} \log \left\| \sigma_A^{(1-\alpha)/2\alpha} \otimes \tau(\varepsilon)_B^{(1-\alpha)/2\alpha} \sigma_{AB}^{1/2} \right\|_{2\alpha} \\
&= \frac{2\alpha}{1 - \alpha} \log \left[\left\| \rho_{AB}^{1/2} \rho_A^{(1-\beta)/2\beta} \otimes \tau(\varepsilon)_B^{(1-\beta)/2\beta} \right\|_{2\beta} \left\| \sigma_A^{(1-\alpha)/2\alpha} \otimes \tau(\varepsilon)_B^{(1-\alpha)/2\alpha} \sigma_{AB}^{1/2} \right\|_{2\alpha} \right] \\
&\geq \frac{2\alpha}{1 - \alpha} \log \left\| \rho_{AB}^{1/2} \left(\rho_A^{(1-\beta)/2\beta} \otimes \tau(\varepsilon)_B^{(1-\beta)/2\beta} \right) \left(\sigma_A^{(1-\alpha)/2\alpha} \otimes \tau(\varepsilon)_B^{(1-\alpha)/2\alpha} \right) \sigma_{AB}^{1/2} \right\|_1 \\
&= \frac{2\alpha}{1 - \alpha} \log \left\| \rho_{AB}^{1/2} \left(\rho_A^{(1-\beta)/2\beta} \sigma_A^{(1-\alpha)/2\alpha} \otimes \tau(\varepsilon)_B^{(1-\beta)/2\beta} \tau(\varepsilon)_B^{(1-\alpha)/2\alpha} \right) \sigma_{AB}^{1/2} \right\|_1 \\
&= \frac{2\alpha}{1 - \alpha} \log \left\| \rho_{AB}^{1/2} \sigma_{AB}^{1/2} \right\|_1 \\
&= \frac{2\alpha}{1 - \alpha} \log F(\rho_{AB}, \sigma_{AB}).
\end{aligned}$$

In the first equality and inequality we used (2.20) and (2.19). The following equalities follow from the definition of the sandwiched Rényi divergence (see Definition 2.4) and (2.22). In the second inequality we apply Hölder's inequality (2.2). For the second-to-last equality we use the fact that $\rho_A = \sigma_A$ by assumption, and that $\tau(\varepsilon)_B$ has full support for $\varepsilon \in (0, 1)$, such that $\tau(\varepsilon)_B^{-1} \tau(\varepsilon)_B = I_B$.

We therefore have

$$\begin{aligned}
\frac{2\alpha}{1 - \alpha} \log F(\rho_{AB}, \sigma_{AB}) &\leq \tilde{D}_\beta(\rho_{AB} \parallel \rho_A \otimes \tau_B) - \tilde{D}_\alpha(\sigma_{AB} \parallel \sigma_A \otimes \tau(\varepsilon)_B) - \log(1 - \varepsilon) \\
&\leq \tilde{D}_\beta(\rho_{AB} \parallel \rho_A \otimes \tau_B) - \min_{\omega_B \in \mathcal{D}(\mathcal{H}_B)} \tilde{D}_\alpha(\sigma_{AB} \parallel \sigma_A \otimes \omega_B) - \log(1 - \varepsilon) \\
&= \tilde{I}_\beta(A; B)_\rho - \tilde{I}_\alpha(A; B)_\sigma - \log(1 - \varepsilon),
\end{aligned}$$

where we chose τ_B as the optimizing state for $\tilde{I}_\beta(A; B)_\rho$ in the last step. Since this relation holds for all $\varepsilon \in (0, 1)$, we obtain the claim by taking the limit $\varepsilon \searrow 0$. \square

Proof of (2.18) in Proposition 2.8. We can rewrite definition (2.14) in Section 2.4 as

$$\tilde{I}_\beta(A; B|C)_\rho = -\frac{2\alpha}{\alpha - 1} \log \left\| \rho_{BC}^{(1-\beta)/2\beta} \rho_C^{(\beta-1)/2\beta} \rho_{AC}^{(1-\beta)/2\beta} \rho_{ABC}^{1/2} \right\|_{2\beta}.$$

Then consider

$$\begin{aligned}
& \frac{1 - \alpha}{2\alpha} \left[\tilde{I}_\beta(A; B|C)_\rho - \tilde{I}_\alpha(A; B|C)_\sigma \right] \\
&= \log \left[\left\| \sigma_{ABC}^{1/2} \rho_{AC}^{(1-\alpha)/2\alpha} \rho_C^{(\alpha-1)/2\alpha} \rho_{BC}^{(1-\alpha)/2\alpha} \right\|_{2\alpha} \left\| \rho_{BC}^{(1-\beta)/2\beta} \rho_C^{(\beta-1)/2\beta} \rho_{AC}^{(1-\beta)/2\beta} \rho_{ABC}^{1/2} \right\|_{2\beta} \right] \\
&\geq \log \left[\left\| \sigma_{ABC}^{1/2} \rho_{AC}^{(1-\alpha)/2\alpha} \rho_C^{(\alpha-1)/2\alpha} \rho_{BC}^{(1-\alpha)/2\alpha} \rho_{BC}^{(1-\beta)/2\beta} \rho_C^{(\beta-1)/2\beta} \rho_{AC}^{(1-\beta)/2\beta} \rho_{ABC}^{1/2} \right\|_1 \right] \\
&= \log \left\| \sigma_{ABC}^{1/2} \rho_{ABC}^{1/2} \right\|_1
\end{aligned}$$

$$= \log F(\rho_{ABC}, \sigma_{ABC}),$$

which yields the claim. \square

B Proofs of Section 3.4: State redistribution with feedback

The following lemma is used to prove the strong converse for state redistribution with feedback, Theorem 3.4. Note that the proof closely follows that of the corresponding result in [8].

Lemma B.1. *Consider the fidelity*

$$F := F\left(\psi_{A'B'C'R} \otimes \Phi_{T_A T_B}^m, (\mathcal{D}_M \circ \mathcal{E}_M \circ \dots \circ \mathcal{D}_1 \circ \mathcal{E}_1)\left(\psi_{ABCR} \otimes \Phi_{T_A T_B}^k\right)\right), \quad (\text{B.1})$$

where the encoding and decoding maps \mathcal{E}_i and \mathcal{D}_i for $i = 1, \dots, M$ are given as in Section 3.4. For $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$, we have the following bounds on F :

$$\log F \leq \frac{1-\alpha}{2\alpha} \left(\log |T_A| - \log |T'_A| + \sum_{i=1}^M \log |Q_i| + \sum_{i=1}^{M-1} \log |Q'_i| - S_\beta(AB)_\psi + S_\alpha(B)_\psi \right), \quad (\text{B.2})$$

$$\log F \leq \frac{1-\alpha}{2\alpha} \left(2 \sum_{i=1}^M \log |Q_i| - \tilde{S}_\beta(R|B)_\psi + \tilde{S}_\alpha(R|AB)_\psi \right), \quad (\text{B.3})$$

$$\log F \leq \frac{1-\alpha}{2\alpha} \left(2 \sum_{i=1}^M \log |Q_i| - \tilde{I}_\alpha(R; AB)_\psi + \tilde{I}_\beta(R; B)_\psi \right). \quad (\text{B.4})$$

Proof. We first prove (B.3). For $\alpha \in (1/2, 1)$ and $\beta = \alpha/(2\alpha - 1)$, we can bound the fidelity F (defined in (B.1)) from above by

$$\begin{aligned} \frac{2\alpha}{1-\alpha} \log F &\leq \frac{2\alpha}{1-\alpha} \log F\left(\psi_{A'B'R} \otimes \pi_{T_B}^m, \sigma_{A'B'RT_B}^M\right) \\ &\leq \tilde{S}_\alpha(R|A'B'T_B)_{\psi \otimes \pi^m} - \tilde{S}_\beta(R|A'B'T_B)_{\sigma^M} \\ &= \tilde{S}_\alpha(R|AB)_\psi - \tilde{S}_\beta(R|A'B'T_B)_{\sigma^M}, \end{aligned} \quad (\text{B.5})$$

where we used the monotonicity of the fidelity under partial trace in the first inequality, eq. (2.16) of Proposition 2.8 in the second inequality, and eq. (2.11) of Lemma 2.7 together with the fact that $\psi_{A'B'R} = \psi_{ABR}$ in the equality. Consider then the following chain of inequalities for the second term on the right-hand side of (B.5) (see Figure 4):

$$\begin{aligned} -\tilde{S}_\beta(R|A'B'T_B)_{\sigma^M} &\leq -\tilde{S}_\beta(R|Q_M B_{M-1})_{\omega^M} \\ &\leq -\tilde{S}_\beta(R|B_{M-1})_{\omega^M} + 2 \log |Q_M| \\ &= -\tilde{S}_\beta(R|B_{M-1})_{\sigma^{M-1}} + 2 \log |Q_M| \\ &\leq -\tilde{S}_\beta(R|Q'_{M-1} B_{M-1})_{\sigma^{M-1}} + 2 \log |Q_M| \\ &\leq -\tilde{S}_\beta(R|Q_{M-1} B_{M-2})_{\omega^{M-1}} + 2 \log |Q_M| \\ &\vdots \\ &\leq -\tilde{S}_\beta(R|BT_B)_{\omega^1} + 2 \sum_{i=1}^M \log |Q_i| \\ &= -\tilde{S}_\beta(R|BT_B)_{\psi \otimes \pi^k} + 2 \sum_{i=1}^M \log |Q_i| \end{aligned}$$

$$= -\tilde{S}_\beta(R|B)_\psi + 2 \sum_{i=1}^M \log |Q_i|. \quad (\text{B.6})$$

In the first inequality we used data processing with respect to the decoding map \mathcal{D}_M (Proposition 2.5(vii)). The second inequality follows from the dimension bound for the Rényi conditional entropy (eq. (2.9) of Lemma 2.7). In the first equality we used the fact that the system B_{M-1} is not affected by the encoding \mathcal{E}_M . The third inequality is data processing for the Rényi conditional entropy with respect to the partial trace over Q'_{M-1} . We then iteratively apply these steps until we reach the last inequality. The subsequent equality follows from the fact that the encoding \mathcal{E}_1 does not act on the systems B and T_B . In the last step we used eq. (2.11) of Lemma 2.7. Combining (B.5) and (B.6) now yields (B.3). The proof of the bound in (B.4) follows in a similar manner, and we therefore omit it.

To prove (B.2), we consider Stinespring isometries $U_{\mathcal{E}_i}$ and $U_{\mathcal{D}_i}$ of the encoding and decoding maps \mathcal{E}_i and \mathcal{D}_i with environments E_i and D_i , respectively. Moreover, in the following calculations we denote by ω^i and σ^i the *pure states* obtained from applying the isometries $U_{\mathcal{E}_i}$ and $U_{\mathcal{D}_i}$ to the initial state $\psi \otimes \Phi^k$, respectively. The final state of the protocol is then the pure state

$$|\sigma_{A'B'C'RT'_A T'_B E_1 \dots E_M D_1 \dots D_M}^M\rangle = (U_{\mathcal{D}_M} U_{\mathcal{E}_M} \dots U_{\mathcal{D}_1} U_{\mathcal{E}_1} \otimes I_R)(|\psi_{ABCR}\rangle \otimes |\Phi_{T_A T_B}^k\rangle).$$

By Uhlmann's theorem there exists a pure state $\chi_{E_1 \dots E_M D_1 \dots D_M}$ such that the following holds for $\alpha \in (1/2, 1)$ and $\beta = \alpha / (2\alpha - 1)$:

$$\begin{aligned} \frac{2\alpha}{1-\alpha} \log F &= \frac{2\alpha}{1-\alpha} \log F \left(\sigma_{A'B'C'RT'_A T'_B E_1 \dots E_M D_1 \dots D_M}^M, \psi_{A'B'C'R} \otimes \Phi_{T'_A T'_B}^m \otimes \chi_{E_1 \dots E_M D_1 \dots D_M} \right) \\ &\leq \frac{2\alpha}{1-\alpha} \log F \left(\sigma_{A'B'T'_B D_1 \dots D_M}^M, \psi_{A'B'} \otimes \pi_{T'_B}^m \otimes \chi_{D_1 \dots D_M} \right) \\ &\leq S_\alpha(A'B'T'_B D_1 \dots D_M)_{\sigma^M} - S_\beta(A'B'T'_B D_1 \dots D_M)_{\psi \otimes \pi^m \otimes \chi} \\ &\leq S_\alpha(A'B'T'_B D_1 \dots D_M)_{\sigma^M} - S_\beta(AB)_\psi - \log |T'_B|, \end{aligned} \quad (\text{B.7})$$

where the first inequality follows from the monotonicity of the fidelity under partial trace, the second inequality follows from eq. (2.15) of Proposition 2.8, and the third inequality follows from Proposition 2.5(ii) and (iii). For the first term of the right-hand side of (B.7), consider the following steps:

$$\begin{aligned} S_\alpha(A'B'T'_B D_1 \dots D_M)_{\sigma^M} &= S_\alpha(Q_M B_{M-1} D_1 \dots D_{M-1})_{\omega^M} \\ &\leq S_\alpha(B_{M-1} D_1 \dots D_{M-1})_{\omega^M} + \log |Q_M| \\ &= S_\alpha(R Q_M C' T'_A E_1 \dots E_M)_{\omega^M} + \log |Q_M| \\ &= S_\alpha(R Q'_{M-1} A_{M-1} E_1 \dots E_{M-1})_{\sigma^{M-1}} + \log |Q_M| \\ &\leq S_\alpha(R A_{M-1} E_1 \dots E_{M-1})_{\sigma^{M-1}} + \log |Q_M| + \log |Q'_{M-1}| \\ &= S_\alpha(Q'_{M-1} B_{M-1} D_1 \dots D_{M-1})_{\sigma^{M-1}} + \log |Q_M| + \log |Q'_{M-1}| \\ &= S_\alpha(Q_{M-1} B_{M-2} D_1 \dots D_{M-2})_{\omega^{M-1}} + \log |Q_M| + \log |Q'_{M-1}| \\ &\vdots \\ &\leq S_\alpha(B T_B)_{\omega^1} + \sum_{i=1}^M \log |Q_i| + \sum_{i=1}^{M-1} \log |Q'_i| \\ &= S_\alpha(B T_B)_{\psi \otimes \pi^k} + \sum_{i=1}^M \log |Q_i| + \sum_{i=1}^{M-1} \log |Q'_i| \end{aligned}$$

$$= S_\alpha(B)_\psi + \log |T_B| + \sum_{i=1}^M \log |Q_i| + \sum_{i=1}^{M-1} \log |Q'_i|. \quad (\text{B.8})$$

In the first equality we used invariance of the Rényi entropy under the isometry $U_{\mathcal{D}_M}$ (Proposition 2.5(iv)). In the first inequality we used subadditivity (Lemma 2.6), and in the second equality we used the duality of the Rényi entropy since $|\omega^M\rangle$ is a pure state. The third equality follows from the invariance of the Rényi entropy under $U_{\mathcal{E}_M}$. We then follow the same steps iteratively, passing from ω^M to σ^{M-1} and ω^{M-1} and so on, until we reach $\omega_{BT_B}^1 = \psi_B \otimes \pi_{T_B}^k$. Substituting (B.8) in (B.7) then yields (B.2), and we are done. \square

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